

University of Minho Informatics Department

Master Course in Informatics Engineering

RAIKON

Invariant Detection Meets Random Test Case Generation

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Abstract

Full fledged verification of software ensures correction to a level that no other technique can reach. However it requires precise and unambiguous specifications of requirements, functionality and technical aspects of the software to be verified. Furthermore, it requires that these specifications together with the produced models and code be checked for conformity. This represents beyond doubt an investment that most developers and companies are neither able nor willing to make.

Although testing can not reach the same level of assurance as full fledged verification, it is the most widely accepted and used technique to validate expectations about software products. Testing is the most natural way of checking that a piece of software is doing what the developers expect it to do. Improvements to test case generation have the potential to produce a great impact in the state of the art of software engineering, by putting Software Testing closer to Formal Software Verification.

This is an exploratory project, aimed at surveying the current state of the art in the field of test case generation and related techniques for the Java language, eventually suggesting possible advancements in the field.

Resumo

A verdadeira verificação de software garante correcção de software a um nível que nenhuma outra técnica consegue igualar. No entanto, exige especificações precisas e inequívocas de requisitos de funcionalidade e aspectos técnicos do software a ser verificado. Além disso, é necessário que as especificações, juntamente com os modelos produzidos e código sejam verificados quanto à sua conformidade. Isto representa indubitavelmente um investimento que a maioria dos profissionais e empresas não são capazes, nem estão dispostos a fazer.

Embora os testes não alcancem o mesmo nível de garantia como a verificação completa, é a técnica mais amplamente aceite e usada para validar as especificações sobre produtos de software. O teste é a forma mais natural de verificar que um pedaço de software cumpre o que os programadores esperam que faça. Melhorias na geração de boletins de teste têm o potencial de produzir um grande impacto no estado da arte da engenharia de software, colocando o teste de software mais perto da Verificação Formal de Software.

Este projecto é de carácter exploratório, visando o levantamento do estado actual da área de geração de casos de teste para a linguagem Java e técnicas relacionadas, sugerindo avanços possíveis nesta área de validação de software.

Contents

Сс	onten	ts																	iii
Li	List of Figures															v			
List of Tables															vii				
Li	st of	Listings																	ix
Li	st of	Acronyms																	xi
1	Intro	oduction																	3
	1.1	Problem S	tatement												_				3
	1.2	Research	Plan																4
	1.3	Backgrour	nd																5
		1.3.1 Tes	tina Philosophi	es															5
		1.3.2 Lev	els of Testina .																6
		1.3.3 Go	als of Testing																6
		134 Ad	equacy Criteria					• •	• •		•••			• •	•	• •		-	7
	1.4	Document	Outline			• •					•••				•				8
2	Lite	rature Rev	ew																9
_	2.1	Review Pr	ocess																9
	2.2	Test Case	Generation Tec	hniques															10
		2.2.1 Ra	ndom Testina																10
		2.2.2 Bo	unded Exhausti	ve Testin	na .														12
		2.2.3 Mu	tation Driven Te	stina															12
		2.2.4 Mo	del Based Testi	na															12
		2.2.5 Su	mmarv																14
	2.3	Auxiliary T	echniques																15
		2.3.1 Co	, nstraint Inferenc	е															15
		2.3.2 Co	nstraint Solving																16
		2.3.3 Inv	ariant Detection																16
		2.3.4 Inp	ut Classification																17
	2.4	Challenge	S																17
		2.4.1 Inp	ut Generation.																18
		2.4.2 Tes	st Oracles																18
		2.4.3 Ge	neral Challenge	S		•••									•			•	18
3	Тоо	I Review																	21
	3.1	Review Pr	ocess	• • • • •													• •		21

CONTENTS

 3.2.1 Test Case generation tools 3.2.2 Interesting Tools for Software Testing 3.3 A Case Study - Checkstyle 3.3.1 Subject of Study: Checkstyle 3.3.2 Tools Tested 3.3.3 Tools Not Tested 3.4 Summary 4 Каком - An Experiment in TCG Tool Interoperability 4.1 Implementation 		· · · · · · · · · · · · · · · · · · ·	· · ·	· · ·	• • • • •	21 37 44 46 49 50
 3.2.2 Interesting Tools for Software Testing 3.3 A Case Study - Checkstyle 3.3.1 Subject of Study: Checkstyle 3.3.2 Tools Tested 3.3.3 Tools Not Tested 3.4 Summary 4 Raikon - An Experiment in TCG Tool Interoperability 4.1 Implementation	· · · ·	· · ·	•	· · ·		37 44 44 46 49 50
 3.3 A Case Study - Checkstyle	· · · ·	· · ·	-	· ·		44 46 49 50
 3.3.1 Subject of Study: Checkstyle	· · · ·	· · · ·		 		44 46 49 50
 3.3.2 Tools Tested 3.3.3 Tools Not Tested 3.4 Summary 4 Raiкоn - An Experiment in TCG Tool Interoperability 4.1 Implementation 		· · · ·		 		46 49 50
3.3.3 Tools Not Tested 3.4 Summary 4 Raikon - An Experiment in TCG Tool Interoperability 4.1 Implementation	· · · ·	· ·			•	49 50
 3.4 Summary 4 Какол - An Experiment in TCG Tool Interoperability 4.1 Implementation 			-		•	50
4 Raiкon - An Experiment in TCG Tool Interoperability 4.1 Implementation						51
4.1 Implementation						21
the improvement of the test of tes						52
4.2 Evaluation Set Up						53
4.3 Results						54
4.3.1 Experiment A: Single Class						54
4.3.2 Experiment B: Package						56
4.3.3 Experiment C: Whole System						56
4.4 Discussion			•		•	56
5 Conclusions						59
Bibliography						65
Index of Terms						73
A Detailed Results						77
B Additional Tooling						99
B 1 BIB2CSV						99
B 2 Fan-In/Fan-Out Calculator	•••	•••	•	• •	•	99

List of Figures

1.1	Typical test case flow.	5
2.1	Concept lattice	11
2.2	Typical workflow of exhaustive testing.	13
3.1	How RANDOOP works [Pacheco and Ernst, 2007].	21
3.2	Software Testing vs. Model Checking [NASA Ames Research Center, 2011]	32
3.3	Example visualisation of a BinaryTree of size 5	39
3.4	Applying ReAssert	41
3.5	AUTOMATIC JUNIT CREATION TOOL main window.	43
4.1	Architecture for Raikon	51
4.2	Flow of the experiments.	54
4.3	Results for FastStack	55
4.4	Results for org.apache.commons.collections.collection	57
4.5	Results for Commons Collections	58

List of Tables

1.1	Cost to fix a defect in various project stages [Kaner et al., 2001]	4
3.1	Evaluation criteria for the case study.	22
3.2	Summary for RANDOOP	24
3.3	Summary for QUICKCHECK.	26
3.4	Summary for JCHECK.	27
3.5	Summary for FEED4JUNIT.	28
3.6	Summary for T2.	30
3.7	Summary for JMLUNITNG.	31
3.8	Summary for Java Pathfinder	33
3.9	Summary for UNITCHECK	35
3.10	Summary for UDITA.	36
3.11	Summary Eclat.	36
3.12	Summary for PEX	37
3.13	Summary for Korat.	40
3.14	Summary for REASSERT.	42
3.15	Summary for DAIKON.	43
3.16	Summary for Automatic JUnit Creation Tool.	44
3.17	Summary for JWalk	45
3.18	Information collected from Checkstyle	46
3.19	Coverage results achieved by the manual test cases	46
3.20	Fan-In and Fan-Out by package of Checkstyle.	47
3.21	RANDOOP generated tests for the target package compared with manually written	
	tests	47
3.22	Comparison of JMLUnitNG results	49
		~ 1
5.1	test case generation (TGG) Tools Summary.	01 00
5.2		52 00
5.3		03
A.1	Set up time for Fast Stack	77
A 2	Set up time for organache commons collections collection	 77
A 3	Set up time for Commons Collections	 77
A 4	Detailed results for Fast Stack (0-55 seconds)	78
Δ 5	Detailed results for Fast Stack (60-120 seconds)	79
A 6	Detailed results for organache commons collections collection (0-55 seconds)	, 0 80
Α 7	Detailed results for org apache commons collections collection (60-120 seconds)	81
A 8	Detailed results for and anache commons collections collection (125-185 seconds)	82
Δ9	Detailed results for ang anache commons collections collection (120 100 seconds) s	83
Δ 10	Detailed results for Commons Collections (0-55 seconds)	84
7.10		54

A.11 Detailed results for Commons Collections (60-120 seconds) .	 85
A.12 Detailed results for Commons Collections (125-185 seconds)	 86
A.13 Detailed results for Commons Collections (190-250 seconds)	 87
A.14 Detailed results for Commons Collections (255-315 seconds)	 88
A.15 Detailed results for Commons Collections (320-380 seconds)	 89
A.16 Detailed results for Commons Collections (385-445 seconds)	 90
A.17 Detailed results for Commons Collections (450-510 seconds)	 91
A.18 Detailed results for Commons Collections (515-575 seconds)	 92
A.19 Detailed results for Commons Collections (580-640 seconds)	 93
A.20 Detailed results for Commons Collections (645-705 seconds)	 94
A.21 Detailed results for Commons Collections (710-770 seconds)	 95
A.22 Detailed results for Commons Collections (775-835 seconds)	 96
A.23 Detailed results for Commons Collections (840-900 seconds)	 97

List of Listings

2.1	Example to illustrate symbolic execution.	15
3.1	Implementation of a random Stack generator	24
3.2	Application of the Stack generator in a test case	25
3.3	Application of the Stack generator in a parameterised test case	26
3.4	Example of a Feed4JUnit equivalence partition test.	28
3.5	Example of T2 specifications.	29
3.6	Example of a test case invoking T2.	29
3.7	Example JPF configuration.	33
3.8	The classic dining philosophers problem.	34
3.9	Example binary tree implementation.	38
3.10	Example binary tree imperative predicate	38
3.11	Example binary tree finitisation.	39
3.12	Two Invariants inferred by DAIKON.	48
3.13	Example of a class annotated with invariants detected by DAIKON.	48
3.14	Example of incomplete invariants detected by DAIKON.	49

List of Acronyms

BBT	•	•	•	-	 •	•	•	•	•	 •	•	•	•	•	•	• •	 •	•	•	•	•	•	•	•	 .Black-box testing
FSV					 •		•	•						•	•			•		•	•				 .Formal software verification
GBT					 		•	•						•	•						•				 .Grey-box testing
PUT					 •		•	•						•	•			•	•	•	•	•	•		 .Parameterised unit testing
TCG					 •		•	•						•	•						•				 .Test case generation
WBT				-	 																				 .White-box testing

Chapter 1

Introduction

Quoting the co-founder of Microsoft, Bill Gates, in a recent interview [Eaton, 2011], the importance of software is higher today than ever. Software is developed by people with no completely reliable method to do so. This means that errors in software, which due to the importance of software, are more and more dangerous.

To verify software there are formal software verification (FSV) techniques, which allow practitioners to achieve the highest degree of confidence in its correction [Holzmann, 2010]. However, these techniques come with overloaded costs. For example, the estimated costs of completely verifying an operating system was reported to be around 20 people years [Klein et al., 2009]. This gives a rough estimate of 700 US\$ per line of code. In the field of verification techniques we can find theorem proving and model checking.

Theorem proving is a formal verification technique that supports program mathematical correction arguments [Chang et al., 1973]. In order to verify a program using theorem provers, practitioners need to establish a sound link between the implementation in a programming language and its mathematical semantics. Only then *off-the-shelf* theorem provers be used to discharge the mathematical logic derived from the properties to be verified [Alexi, 1988; Vermolen, 2007]. Model checking is another technique to formally verify software. It is based on an exhaustive search for counter examples for a given property, within a bounded domain [Clarke and Emerson, 1981; Clarke, 2008]. Compared to theorem proving, it is easier to automate, but does not achieve the same degree of completeness, due to its bounded nature.

Neither of these techniques, reliable as they are, is widely adopted in mainstream software industry [Holzmann, 2010]. This is due to the cost/benefit trade-off in applying them [Klein et al., 2009]. The alternative that is most accepted is testing. Testing is even lighter than model checking [Myers, 1979], at the cost of being even less complete. What makes software testing more affordable is the fact that it can be done by regular developers using the same programming languages in which products are implemented. Testing relates inputs with outputs in accordance with an expected results [Myers, 1979].

1.1 Problem Statement

A good testing infrastructure for a product or service, it is estimated at least 50% of the total costs of the project [Myers, 1979]. This is due to the importance of detecting problems in the early stages of a project, when they are cheaper to fix [Sommerville, 2010]. An estimation of the costs required to fix problems detected in earlier product development stages can be found in Table 1.1 [Kaner et al., 2001]. Note however that, there is also a cost associated with the maintenance of the test suites throughout the entire project, to ensure the correct behaviour of the product throughout the various stages [Harrold et al., 1993]. An inadequate testing infrastructure can lead to even bigger losses [Tassey, 2002].

CHAPTER 1. INTRODUCTION

	Stage Detected												
Stage Introduced	Requirements	Architecture	Construction	System Test	Post-release								
Requirements	1x	Зx	5x	10x	10-100x								
Architecture	-	1x	10x	15x	25-100x								
Construction	-	-	1x	10x	10-25x								

Table 1.1: Cost to fix a defect in various project stages [Kaner et al., 2001]

In order to reduce costs, or to better take advantage of them, researchers and practitioners focus their efforts in ways to automate test case generation (TCG) [Edvardsson, 1999]. This technique attempts to cover a considerable amount of inputs, while also testing more production code. TCG attempts to increase the effectiveness of testing, since it is less exposed to trivial mistakes. Also, due to the generation being automatic, there is considerably less effort with maintenance of a test suite. On the downside, automatically generated tests tend to be simpler than manually written ones, and therefore they are only able to detect simpler flaws in the software.

1.2 Research Plan

The ultimate goal of this research is to assess the prospect of advance in TCG and, if possible, provide an implementation of such advance. First we need to explore what techniques are available. Only then we will know whether or not potential for making a contribution to the field. For the purposes of this research, we will focus on TCG for the Java programming language. However, we will not limit the study of techniques and tools to the ones for Java, as there might be valuable lessons to learn from other languages or frameworks.

We propose to answer the following research questions:

- RQ1 What methods are there for test case generation?
- RQ2 How do the existing methods compare to each other?
- **RQ3** Is there an opportunity to produce an advancement in existing methods for fully automated test case generation?

In order to answer these research questions we have decided to break each of them into tasks in order to better structure the research. **RQ1** is broken down into the following tasks:

- **T1.1** Research the state of the art in TCG;
- T1.2 Research techniques that might be used in the context of TCG;

For RQ2 the break down in tasks is:

- T2.1 Select of the most promising tools according to the survey done to answer RQ1;
- T2.2 Conduct of a case study, where the tools selected in T2.1 are applied to real world software;

Finally RQ3 is broken down into the following tasks:

- **T3.1** Formulate requirements for our contribution;
- **T3.2** Based on the results of the case study conducted for **RQ2**, decide whether to improve one of the existing tools or implement a new tool chain;
- T3.3 Implement a prototype of what was decided in T3.2;

1.3 Background

A test case consists of passing an input value to a method under test and checking if the output of that method conforms with a set of expected results in order to confirm that the method is behaving correctly [Myers, 1979]. Figure 1.1 illustrates a typical flow of a test case. The main advantage of manually written test cases is the control it provides to developers, allowing them to define input and expected output values. The main disadvantages are the high costs of the process and the obvious subject to human error.



Figure 1.1: Typical test case flow.

The following sections provide a deeper understanding about testing.

1.3.1 Testing Philosophies

By tradition, testing philosophies are divided into two major factions, white-box testing (WBT) and black-box testing (BBT) [Myers, 1979]. Their combination leads to get another philosophy known as grey-box testing (GBT) [Kaner et al., 1999].

WBT is a testing methodology that checks the internal workings of a piece of software. To follow this approach the source code must be available to the practitioner, since test data are derived from the logic of the software. This approach is useful in detecting errors and problems within the code itself, but fails to detect unimplemented or incorrect specifications, because WBT completely disregards specifications. Ultimately, the goal of WBT is to check every possible execution of the software [Myers, 1979].

On the other hand, BBT validates that a software conforms to a given specification. Unlike WBT, tests following BBT philosophy do not require knowledge of the source code. This approach is used to detect incorrectly implemented specifications. However, it is not as accurate as WBT in pin-pointing problems with the internal logic. In the long run BBT attempts to test all possible input/output combinations [Myers, 1979].

Blending some aspects of both philosophies has lead to a third one, known as GBT. GBT does not fit into BBT because it leverages some knowledge of the internal logic. However, it can not be considered WBT either since it does not aim at exploring the problems in the internal logic of the software. The main idea is to leverage knowledge of the specification and the internal logic to produce better informed tests cases [Kaner et al., 1999].

CHAPTER 1. INTRODUCTION

Summing up, these philosophies offer guidelines to create reliable test suites. Neither BBT nor WBT, present a better offer than the other. On the other hand, GBT offers a good compromise between specification validation and structural verification.

1.3.2 Levels of Testing

In the process of testing software products, there are various *granularity* levels on which practitioners can test software. The finest of them is unit testing, which leads to integration testing, and finally to the coarsest level, system testing [Kaner et al., 1999]. Other levels exist, however, these three are considered the most relevant. This section will describe each of them in detail.

The finest level, unit testing, consists in testing software units, the smallest components of a software system [Kaner et al., 1999]. In essence a unit test is a call to the unit being tested with a predefined input value. By carefully selecting the inputs passed to the unit under test, the tester can predict the expected outputs, and then express that expectation as an assert statement. Such statements in a test are called oracles. In order to write high quality unit tests, both a very rigorous discipline and considerable effort are required, since each unit requires at least one test, but preferably one test per control path [Zhu et al., 1997].

There is, a technique that attempts to reduce efforts in unit testing, known as parameterised unit testing (PUT) [de Halleux and Tillmann, 2008; Tillmann et al., 2010]. The major contribution introduced by this technique is the possibility of providing parameterised input values, meaning that input values are not defined in the test itself, but at run-time and according to a predicate expression. PUT is therefore a generalisation of unit testing that requires more abstract oracles. Both approaches follow a WBT philosophy.

A coarser level of testing is integration testing. At this level, instead of testing single units, entire modules or sets of units are tested [Myers, 1979; Kaner et al., 1999]. The main goal is to check that the interactions of the modules under test are behaving as expected. However, there are some limitations to integration testing, for instance it makes it harder to pin-point the source of the detected problems [Kaner et al., 1999]. Integration testing usually follows a GBT philosophy.

The coarsest level of testing is called system testing [Kaner et al., 1999]. At this level the whole product is placed under test. This is done to validate that the software meets the requirements defined by the stakeholders, which makes system testing purely BBT. System testing suffers from the same limitations as integration testing, making it even harder to pin-point the source of flaws [Kaner et al., 1999].

Summing up, three of the most relevant levels of testing were described: unit testing, integration testing and system testing. No level stands out from the others, meaning they all complement each other. While testing a software system, one should consider more than one level of testing to ensure a higher degree of confidence [Kanstrén, 2008].

1.3.3 Goals of Testing

What do testers expect to learn from testing a piece of software? Testers may want to learn more about the performance of the system under test or if all requirements are met. Approaches known as non-functional testing, regression testing and acceptance testing provide for different goals in testing.

Non-functional testing addresses non-functional requirements, such as performance and scalability [Myers, 1979]. To evaluate performance of the software, the system is run several times in different machines and the performance on each machine is measured. To measure scalability, the system is put through an increasing amount of workloads in order to assess the workload which it can handle.

Regression testing is a type of testing whose aim is to find errors that might have been introduced after applying changes to a certain functionality [Kaner et al., 1999]. This is done by recording in a regression test suite the behaviour of software before the changes are made. Then, after making the changes the regression test is used to check that only the intended new behaviour was affected by the change. Thus, no undesirable side effects were introduced. Regression tests may reveal newly introduced faults, or they may fail if the changes were meant to intentionally change functionality. Another use of regression testing, other than finding errors, is to assert that the performance of the system was not affected. Regression testing can, therefore, be accomplished at the various levels of testing.

On the other hand there is acceptance testing, also known as functional testing, which is a type of testing to check that a given implementation meets its requirements [Kaner et al., 1999]. This can be done by a tester with access to the requirements of the system, or by the final users of the system. This is usually a final stage of development, where any flaws may be uncovered before delivery. Usually, acceptance testing is done at system testing, but it is also possible to apply at the other levels.

These two forms of testing are not necessarily opposites, as they complement each other. An acceptance test can be later used as a regression test. While regression tests are based on previous versions of the production code, acceptance tests are based on a baseline specification that represents the intended functionality [Myers, 1979].

1.3.4 Adequacy Criteria

As previously mentioned, the main goal of software testing is to provide stakeholders with confidence in the software under test. In order to be easily understood by everyone, confidence should be expressed as a number. As such, the quality of a given test suite must be measured according to proper adequacy criteria [Zhu et al., 1997]. This section will describe some adequacy criterion: structural coverage; fault detection; input bound coverage and assertion density.

The first adequacy criterion covered in this section is also the most common, structural coverage. This criterion measures how much of the source code is executed by the test suite. Structural coverage can measure the coverage over structural elements like: line coverage, branch coverage and loop coverage [Zhu et al., 1997].

It is believed that when an element is executed by a test suite there are less chances of problems occurring in that element [Zhu et al., 1997]. This is, however, not enough because in addiction to executing the statement, it is also necessary to check that the result is the expected.

Another criterion of relevance is fault detection. This criterion measures how reliable a given test suite is in detecting defects in the code [Zhu et al., 1997]. This becomes especially useful while creating a test suite for regression testing. These tests aim to prevent the introduction of faults while maintaining the existing features and in developing new ones. A common form of evaluating the robustness of a test suite in relation to fault detection is *mutation analysis* [Zhu et al., 1997]. This analysis consists in injecting artificial changes, *mutations*, in the source code and then checking how many of the *mutations* the test suite is able to detect. The higher the number of *mutations* detected the more reliable a given suite is.

Another criterion is input bound coverage. This adequacy criterion measures how much of the input domain is being covered [Zhu et al., 1997]. The higher the input bound coverage, the more

CHAPTER 1. INTRODUCTION

likely the chances of detecting errors in the behaviour of the system under test are. However, covering the whole extent of the input domain is a virtually impossible task in most cases.

The last adequacy criterion covered in this section is known as assertion density. This criterion measures the number of assertions per lines of code [Foster et al., 2003], thus the higher the ratio the more assertions per each line of code. This metric leads to a higher confidence in the test suite, in the sense that more assertions can detect more discrete problems in the software.

In summary, this section provides some insight to several adequacy criteria in the field of software testing. These criteria are very important, because in order to measure the confidence in a software through its tests on should have confidence in the test suite itself.

1.4 Document Outline

The remainder of this dissertation is organised as follows. A review over TCG literature is presented in Chapter 2: Literature Review which also includes an overview of the current major challenges in TCG which allow us to answer **RQ1**. Chapter 3: Tool Review describes a review of chosen tools that can be used in TCG, we try to apply all the selected tools to a real world tool, Checkstyle.. Chapter 4: RAIKON - An Experiment in TCG Tool Interoperability describes the implementation of our tool, Raiko, which is a combination of a tool for emphdirected random testing, RANDOOP, with a constraint inference system, DAIKON. The document concludes in Chapter 5: Conclusions drawing conclusions and future work.

Chapter 2

Literature Review

This chapter reports on the literature review process conducted while researching into the state of the art in TCG. The most relevant TCG techniques found in the review are described, including random testing, bounded exhaustive testing, mutation-driven testing and finally model-based testing. Other techniques that can be used in the context of TCG are also discussed. The chapter closes with a discussion of the identified challenges in test case generation.

2.1 Review Process

This section will describe the literature review process conducted while studying the state of the art in software testing. The papers were found through searches in known search engines and through analysis of bibliography listings. It was decided to experiment with formal concept analysis [Wolff, 1993] in the context of literature reviewing. formal concept analysis provided a wider perspective over the field of TCG, while revealing interesting relations among the many contributions to the field.

Formal concept analysis is a technique that allows the automatic derivation of an ontology from a matrix of objects and attributes [Wolff, 1993]. For the purpose of this review, the objects are the papers and the attributes are the topics covered by the papers. A list of the attributes chosen follows.

Adequacy Criterion If the paper discusses relevant adequacy criteria for test suites;

Constraint Inference If the paper discusses constraint generation techniques;

Constraint Solving If the paper discusses constraint solving techniques;

Data Generation If the paper discusses or addresses the test data generation problem;

Dynamic Analysis If the paper discusses applications of concrete executions;

Path Selection If the paper discusses path selection decidability problems;

Programming Language If the paper relates to a specific programming language;

Software Failure If the paper discusses or describes software failures or faults;

Software Testing If the paper discusses or addresses software testing in some way;

Symbolic Execution If the paper discusses applications of symbolic execution;

TCG Technique If the paper describes a TCG technique;

TCG Problem If the paper describes the TCG problem;

TCG Tool If the paper describes a tool or system that either performs TCG or can be reused as part of a TCG tool-chain.

Having constructed the context matrix, the CONEXP tool [Tilley, 2004; Yevtushenko, 2011] was used to generate the ontology illustrated in Figure 2.1. This ontology is represented by a concept lattice. The bottom node of the lattice collects all papers that do not fit into any of the attributes. By contrast, the top node contains all papers that fit into all of the attributes. However, in the lattice of Figure 2.1, there are no papers neither in the bottom nor top nodes of the lattice. This means that all the papers reviewed were both classifiable and diverse. By examining the remaining nodes, some conclusions can be drawn about the literature review thus carried out. Some of these conclusions follow:

- Most researched techniques for TCG refer to an implementation. This means that all techniques researched have associated tools.
- Not all papers are related to testing. This is due to the need of collecting some information, not directly related with software testing. For example, applications of given techniques in contexts other than software testing.

This experience with formal concept analysis was a pleasant experience. We have found that we had overlapping attributes which lead us to reduce them in number. It has allowed us to acquire a better perspective on what we were searching for and even find some related concepts. Bellow we fine a detailed account of what we have learned from the literature review process.

2.2 Test Case Generation Techniques

This section covers the most common techniques for TCG. These are random testing, bounded exhaustive testing, mutation-driven testing and finally model-based testing.

2.2.1 Random Testing

The most basic TCG technique is random testing. This technique ranks among the most simple techniques to implement, and also as one able to generate a high number of test cases per unit of time [Edvardsson, 2002]. This is due to the simple nature of this technique, since there is no sort of *guidance* in the generation process, i.e. it generates inputs regardless of the domain of the software, as any random stream of data can be used to produce an input. There are essentially three types of application of random testing. It can be used to generate random sequences of method invocations and random input values to execute them, as happens in RANDOOP [Pacheco and Ernst, 2007]. Another possible application of random testing is QUICKCHECK [Claessen and Hughes, 2000; Hughes, 2009].

As a downside to random testing, the chance of achieving an acceptable level of structural coverage is low [Deason et al., 1991; Offutt and Hayes, 1996]. This is due to the nature of the random test generator, which might not be able to produce the required input to cover relevant paths of a method. Recent research, however, has revealed some interesting properties about random testing [Ciupa et al., 2007]. For instance, the number of detected flaws seems to be



Figure 2.1: Concept lattice

CHAPTER 2. LITERATURE REVIEW

inversely proportional to the elapsed testing time, meaning that there is no added gain in increasing the generation time of a random testing tool indefinitely. The same study has also reported that random testing is not limited to detecting artificial flaws seeded into the source code. It showed that this technique is also capable of detecting complex bugs in software.

2.2.2 Bounded Exhaustive Testing

Another TCG technique that is being brought forward is bounded exhaustive testing [Sullivan et al., 2004]. This technique generates test suites to meet a considerably high input bound coverage. The technique ensures that every input combination within a bounded scope is tested to check that the method under test is producing the expected outputs for each and every single one of them, making it a very effective technique for testing. There are two ways of using bounded exhaustive testing: 1) provide an oracle to be checked for a given input bound; 2) provide a description of valid inputs and generate appropriate inputs from the provided description, and pass them to a PUT.

There are two interesting tools in the field of bounded exhaustive testing, namely KORAT [Milicevic et al., 2007] and JAVA PATHFINDER [Visser and Mehlitz, 2005; Pasareanu and Rungta, 2010]. KORAT allows users to describe valid inputs of complex structures and a scope, upon which it generates proper inputs. JAVA PATHFINDER allows users to provide a test oracle. It will check that the software is meeting them. Nonetheless, there are other tools that could be considered for bounded exhaustive testing, even though they are not directed at this. Such is the case of QUICKCHECK [Claessen and Hughes, 2000; Hughes, 2009] with the aid of custom-made data generators for such a purpose [Runciman et al., 2008]. Figure 2.2 illustrates the typical workflow of a bounded exhaustive testing tool.

There are, however, serious drawbacks to bounded exhaustive testing. One of these is the large amount of computing resources required to generate sufficient inputs to cover a significant portion of the input domain, which, in many cases, is virtually unlimited. Nonetheless, recent research advances are attempting to reduce such costs by some orders of magnitude [Jagannath et al., 2009].

2.2.3 Mutation Driven Testing

Mutation-driven testing [Fraser and Zeller, 2010] consists in improving test suites by generating tests to detect a higher number of *mutations*. This technique uses *mutation analysis* techniques (recall in Section 1.3) to learn of possible *mutations* and then it attempts to generate appropriate oracles to properly detect those *mutations*. A mutation-driven testing technique which was implemented in a tool named $\mu Test$ [Fraser and Zeller, 2010] generates test cases while leveraging genetic algorithms to do so. This technique depends on the source code and an existing test suite.

2.2.4 Model Based Testing

Another research approach to TCG is model-based testing. This technique consists in extracting test cases from an existing model of the system under test [Bruno et al., 1995]. There are various approaches to model-based testing, from which, the following are singled out [Utting and Legeard, 2006]:

· Generation of test input data from a domain model;



Figure 2.2: Typical workflow of exhaustive testing.

- · Generation of test cases from an environment model;
- · Generation of test cases with oracles from a behaviour model; and
- · Generation of test scripts from abstract tests.

The first approach, input generation from a model consists in extracting information about the input domain from the model and generate appropriate test data that can be used while testing the software. The second approach, generation of test cases from an environment model where the system under test will be used, for instance a statistical study of the expected load of the system under test. From this model it is possible to construct sequences of calls to the system under test to validate the behaviour under the designated environment. The third approach to model-based testing, and perhaps the most interesting, allows the generation of test oracles, once a behavioural model is available. This type of model holds enough information to generate complete test cases with proper test oracles. Unlike the other approaches, this approach covers all the test construction phases since it contains the relationship between the inputs and the outputs. The last approach aims at the generation of executable test cases based on abstract representation, for example a UML sequence diagram.

The most serious limitation in applying model-based testing is the constant lack of appropriate models to generate tests from. But even in the presence of models, there are still limitations to model-based testing [Utting and Legeard, 2006]. For instance, no approach is able to detect all the differences between the model and the implementation, or the need to prove that the model is correct. Its costs in comparison to other TCG techniques are still relatively high, since it requires abstract models. Also there is the problem of keeping the software model synchronized with the software code. These limitations keep most mainstream software industries away from model-based testing [Sarma et al., 2010].

2.2.5 Summary

This concludes the state of art in TCG. The techniques surveyed were: random testing, bounded exhaustive testing, mutation-driven testing and model-based testing. Random testing is the simplest and cheapest technique studied. However, it is the one with the least chances of finding software bugs. This does not mean that random testing is worthless since there are reports of its use-fulness [Ciupa et al., 2007]. Bounded exhaustive testing has the goal of covering a significant portion of the input domain, increasing the chances of finding a bug on the system under test. Mutation-driven testing is a technique that focuses on the generation of test cases specifically aimed at *mutation* detection. The last TCG studied, model-based testing, leverages a model of the software and generates complete test cases or appropriate input test data to check that the implementation meets the behaviour described by the model of the software.

In summary, no single technique presents a major advantage over the others and they rather complement each other. However, the technique with most potential for TCG is model-based testing, because it is capable of automatically generate the most complete input test cases with test data and test oracles. Also because the generated test cases are capable of validating the implementation against the model, in spite of the constant lack of suitable models for model-based testing preventing this technique from being more widely adopted.

2.3 Auxiliary Techniques

The literature review revealed software testing is commonly related to constraint satisfaction problems. This happens because constraint satisfaction problems are related to the generation of inputs and the inference of oracles, which are two important aspects of TCG.

This section explains constraint inference, constraint solving and related techniques, as well as their applications in the context of TCG.

2.3.1 Constraint Inference

Inference is the process of deriving strict logical consequences from assumed premises. A possible definition of constraint inference is the relation between constraints and their consequences. For example let D and C be two constraint sets. We say that $D \models C$ if and only if, for a given valuation¹ V, if V satisfies D then V also satisfies C. In (2.1) we attempt to express the definition of constraint inference as a mathematical equation.

$$D \models C \iff \forall V : (V(D) \vdash \top \Longrightarrow V(C) \vdash \top)$$
(2.1)

In other words, constraint inference attempts to model a property system into a simpler one. This is done so as to better understand the constraint system and to facilitate its solving. In computer science, constraint inference is used in many fields, like symbolic execution [Ramamoorthy et al., 1976], dynamic analysis [Gulwani et al., 2009] and more recently a combination of both, named concolic execution [Pasareanu and Visser, 2009].

Symbolic Execution

Symbolic execution is a technique that executes programs using symbolic values, this is, instead of using actual values to execute the program, symbolic execution uses predicate expressions that represent the inputs. These expressions act like variable substitutions. The goal of symbolically executing a program is to construct an expression based on its input variables [Ramamoorthy et al., 1976]. This expression can then be solved to compute inputs for the program. Consider the statements in Listing 2.1. A symbolic execution on the statements would produce the expression (((a * a) + (b * b))) = (c * c).

```
x = a * a;
y = b * b;
z = c * c;
result = ((x + y) == z);
```

Listing 2.1: Example to illustrate symbolic execution.

Using symbolic execution, there is no need to perform violation checks on each branch, since all of them can be resolved in a single execution [Ramamoorthy et al., 1976]. However, symbolic execution is entailed by some limitations, for instance function calls can not be resolved, unless we either perform an inline substitution of the function being called for its corresponding instructions or learn the symbolic expression of those functions before finishing the current function. Other

¹The assignment of values

CHAPTER 2. LITERATURE REVIEW

limitations of symbolic execution, are object, pointer and array referencing [Ramamoorthy et al., 1976]. Symbolic execution has no way of knowing their value since they are only instantiated during run-time.

Dynamic Analysis

Dynamic analysis, on the other hand, makes use of actual values while executing a program [Gulwani et al., 2009]. By using actual values to perform the executions, the problem of object, pointer and array referencing is much less accentuated since the actual value of those elements is always known throughout the execution of the program. Also, function calls do not present as much of a problem to dynamic analysis as compared to symbolic execution, since it can simply execute that call in order to know what the returned value is. There is a downside, however, and this is that dynamic analysis requires a considerably large number of executions in order to infer a reliable constraint system. This is due to the violation checks for each branch, which requires different executions, with different input values.

Concolic Execution

Recent research, has lead to a combination of both strategies, symbolic execution and dynamic analysis, leveraging the stronger points of each technique to balance each other weaknesses. This means that the ability to solve branch expressions from symbolic execution is used to reduce the number of executions needed for dynamic analysis. While dynamic analysis can provide values for referenced elements like objects, pointers and arrays, thus eliminating the referencing limitation of symbolic execution. This technique is known as concolic execution [Pasareanu and Visser, 2009]. The name is a contraction of concrete and symbolic execution.

2.3.2 Constraint Solving

Constraint solving is the process of attempting to find a solution to a constraint system. There are some constraint solving techniques that are used inside a TCG. These techniques are used by a TCG to search for solutions to the constraint system representing either the system under test or appropriate inputs to the system under test. Solutions to the system normally guide the TCG towards a desired adequacy criteria. The most relevant technique is known as satisfiability solving [Dixon et al., 2011]. An approach to solving a satisfiability system is conflict-driven [Dixon et al., 2011; Prasad et al., 2005]. This approach searches for a contradiction of the formula to satisfy and if no contradicting example is found, then the formula must be true. Another approach to solving a satisfiability system is based on stochastic local search methods [Prasad et al., 2005], but they do not perform well on structured instances obtained from real verification problems.

2.3.3 Invariant Detection

The previously discussed constraint inference techniques allow for the construction of constraint systems. These techniques have both direct and indirect applications to TCG. Direct applications include the construction of a constraint system that represents the system under test. This constraint system then can be solved to provide guidance towards given adequacy criteria. Indirect applications of constraint inference include invariant detection, a technique that attempts to infer

function invariants. There it more than one way to accomplish this, be it through specification mining [Ammons et al., 2002; Ernst et al., 2007] or static observation of source code [Ramanathan et al., 2007].

Specification mining attempts to dynamically infer program specifications from execution traces of the program [Ammons et al., 2002]. DAIKON [Ernst et al., 2007] is a tool that implements specification mining. It takes as input executions provided by a test suite and it observes properties that were kept along the execution. This leads to a very interesting approach to testing, since test cases can be created from specifications and specifications can be inferred from the execution traces provided by test cases. This creates a feedback loop [Xie and Notkin, 2003], where tests can be generated from specifications and specifications can be refined from tests. However, not all test cases contribute in the same way to this process [Dallmeier et al., 2010]. The more complete the test cases the more reliable the inferred properties will be.

There are also some applications of invariant detection using static analysis [Ramanathan et al., 2007]. By leveraging data flow analysis to gather predicates for each program point, it is possible to derive valid properties that should be valid before and after the execution of a method.

Nonetheless, automatically inferred specifications are not completely trustworthy. Studies reported in [Polikarpova et al., 2009] were made to compare programmer-written and automatically inferred invariants. These studies conclude that automatically inferred contracts do not always cover as much of the program behaviour as the ones manually defined by programmers. In some cases, automatically inferred contracts encounter unspecified or unexpected program behaviour that the developer missed. However, these studies suggest that the automatically inferred contracts could be used either to complete programmer-written contracts or to serve as a baseline for them.

2.3.4 Input Classification

Other relevant techniques that can aid in the process in testing, especially in test data generation, are known as input classification techniques [Carbin and Rinard, 2010; Pacheco and Ernst, 2005]. These techniques consist in classifying inputs based on their potential to reveal faults in the target software.

A possible approach to input classification is to use a model of the operation of the system under test to generate random inputs and then executing the target software with the generated input, checking its behaviour according to the previously extracted model of operation of the software. If the execution behaves incorrectly then it is possible that either an illegal input was found or a possible fault in the software was detected. This approach has been implemented in EcLAT [Pacheco and Ernst, 2005], a tool that generates fault revealing test cases, given an example execution of the system under test.

2.4 Challenges

Software testing and TCG remain a challenging field in software sciences. There are four major trends in software testing: Universal Test Theory, Test-Based Modelling, Fully Automated Testing and Efficacy-maximised Test Engineering [Bertolino, 2007]. To complete this chapter we discuss some of the most relevant challenges remaining in such trends.

2.4.1 Input Generation

The problem in input generation is the difficulty of generating appropriate inputs, i.e. that meet the entry condition for a given method. Some believe that a minimalist set of good program inputs covering a significant amount of branches can reveal the same faults as a huge test data set would [Carbin and Rinard, 2010; Pacheco and Ernst, 2005; Milicevic et al., 2007]. Also the automatic generation of test data is seen as an important step to achieve fully automated testing.

The areas expected to produce considerable advances in input generation are model-based approaches [Bruno et al., 1995], guided applications of random data generators [Pacheco and Ernst, 2007], and a wide variety of search-based techniques [Maragathavalli, 2011]. From model-based approaches it is possible to to use models as a baseline to derive complete test cases [Bertolino, 2007]. This approach could be used in conjunction with concolic execution to guide the execution process and thus traversing the most relevant paths to validate the model. Model-based approaches remain the most promising in their allowing developers to validate their implementation against application models defined by the stakeholders, assuming the model is correct.

Random generation techniques have been put aside in favour of other, more systematic, techniques that were deemed more comprehensive and capable of achieving much better results, like model-based testing and bounded exhaustive testing. Recently, however, random generation have been compared to such systematic techniques and shown to achieve almost identical results with much less effort [Ciupa et al., 2007]. For this reason, researchers currently believe that it is possible to somehow enhance random generation techniques by making them smarter [Pacheco et al., 2007], by leveraging previous executions to better predict the input bound.

2.4.2 Test Oracles

Another challenge in software testing and TCG in particular is the availability of test oracles [Staats et al., 2011; Bertolino, 2007]. Recall that the part testing that decides whether test cases succeed or fail is the oracle. There are several ways to produce test oracles. The most reliable ones involve manual interaction from the developer and even these are mostly incomplete, resulting in test oracles that do not cover all the behaviour they should. The current challenges is to improve current means of producing reliable test oracles, and of course to take test automation a step further by fully automating oracle generation. This is a challenge because there is no way of knowing what to expect from a method with incomplete information. It is only possible to capture the current result and assert that it is always the same. Even then it is necessary to convert observed behaviour into expressions to be asserted. Model-based testing and executable specification languages like JML are regarded as the most promising techniques capable of solving this challenge [Baresi and Young, 2001; Bertolino, 2007].

2.4.3 General Challenges

The so-called general challenges in software testing go beyond the technical aspects of the testing process. General challenges include enrichment to the whole theory behind software testing, like the assessment of the true effectiveness of testing and also how can the results of a test case could influence another test case, or even how the order of which test cases are executed influences the testing process. Also it is very important to know the real costs of testing and the real costs of not testing software applications.

Test effectiveness is related to adequacy criteria. It is considered of vital importance to determine which adequacy criteria or which combination of criteria is best in order to detect likely faults while minimising the costs and effort of testing at the same time. Currently there is no single criterion that stands out from the others, the general consensus being that a combination of criteria is considered best [Bertolino, 2007].

Test composition relates to the order by which test cases are executed. When a given test suite gets too big to run in a single execution, usually the approach is to apply the old divide and conquer - *divide et impera* - strategy [Bertolino, 2007]. However, it has been suggested that by properly ordering test cases it is possible to minimise the effort to run the test suite. Basically the results from previously executed test cases can be used to reduce the effort in executing the whole test suite.

Another interesting challenge in software testing is the inference of *Test Patterns*. Such patterns would be analogous to the well known *Design Patterns* [Gamma et al., 1995] for software. The purpose of these patterns would be to help in detecting the best strategy to thoroughly test a given component or system, thus allowing testing to go faster and simpler. It would also provide for reusable solutions to common problems.

Current state of the art cost evaluation of software testing is only based on estimations. What is also estimated is the cost of poorly testing or not testing applications [Tassey, 2002]. Knowing the real costs and risks of whether testing and not testing software is vital for the future of the software testing discipline as it will allow the assessment between cost and risk for each component and prioritise testing effort.

Finally, a challenge that will always prevail is the proper education of software testers [Bertolino, 2007]. Even with a wide array of tools and techniques available the human eyeball should always have the final say in testing, and good educators in the software testing will make them more aware of common problems and how to check for them.
Chapter 3

Tool Review

The motivation behind the tool review reported in this chapter is to find out what tools are available that either preform TCG, or aid in a testing environment. This chapter describes the review process, the tools reviewed and their range of application. Finally this chapter presents some considerations about the tools.

3.1 Review Process

To learn more about TCG methods some tools were reviewed and analysed under a given set of criteria in order to find out their usefulness in a software development environment. The evaluation criteria used throughout the review are presented in Table 3.1 as well as a summary of what each criterion means and what values it might take throughout the review.

The tools reviewed were found throughout the search for literature in the field, through searches in known search engines and discussion with peers. The tools are mostly focused around the Java programming language, with few a exceptions.

3.2 Reviewed Tools

3.2.1 Test Case generation tools

This section describes the most promising tools for TCG.

Randoop



Figure 3.1: How RANDOOP works [Pacheco and Ernst, 2007].

Language	Indicates the language for which the tool is intended.
Price	Indicates the price or licensing limitations of the tool.
Level of Tests	Indicates the level of tests in which the tool operates, it may take the values Unit, Integration or System testing.
Scope	Indicates the type of the tool, framework, test case generator, etc.
Adequacy Criterion	Indicates the adequacy criteria that the tools attempts to meet.
Incremental	Indicates if the tool leverages any existing test suite.
Previously Generated Tests	Indicates the integration of the tool with previous generated tests.
Integration With Manual Tests	Indicates the level of integration between the tool and manually written tests.
Integration With Development Tools	Indicates the integration of the tool with other development tools.
Code or Binary Level	Indicates the level where the tool operates.
Configuration Effort	Indicates the level of effort required to use the tool: low, medium, high.
Quality of Produced Tests	Describes the quality of the generated tests.

Table 3.1: Evaluation criteria for the case study.

RANDOOP was developed by Carlos Pacheco and Michael Ernst [Pacheco and Ernst, 2007; Pacheco, 2009]. This tool is capable of generating regression testing and contract violation tests, with test inputs and oracles, for Java and .NET¹ software. The underlying technique in RANDOOP is feedback-directed random TCG [Pacheco et al., 2007]. To put it simply, RANDOOP generates sequences of method and constructor invocations. These sequences are built incrementally by randomly selecting which method to apply next to an already existing sequence. Once the sequence is built it is executed in order to determine if it is redundant, illegal, contract-violating, or useful for generating more inputs. When a method invocation requires input parameters, RAN-DOOP randomly selects suitable inputs from a pool of previous invocations or from a pre-defined set of values. Oracles are generated by observing the object state of the software under test and generating proper assertions that check that object state. Figure 3.1 [Pacheco and Ernst, 2007] attempts to illustrate this behaviour. RANDOOP has been validated using a considerable number of targets [Pacheco, 2009]. It showed potential as a reliable tool in generating and finding regression errors and contract violations. Do note that the figure presented may be misleading. Even though there are plans to allow user-defined contracts in the future, currently RANDOOP only checks for a handful of hard-coded contracts, as listed below:

- Equals to null: o.equals(null) should return false;
- Reflexivity of equality: o.equals(o) should return true;
- Symmetry of equality: o1.equals(o2) implies o2.equals(o1);
- Equals-hashcode: If o1.equals(o2)==true, then o1.hashCode() == o2.hashCode()
- No null pointer exceptions: No NullPointerException is thrown if no null inputs are used in a test.

QUICKCHECK for Java

QUICKCHECK was originally a combinator library for the Haskell programming language [Claessen and Hughes, 2000]. It was later implemented for several languages like Java [Open Source Contributors, 2011b,a], Erlang [Hughes and Arts., 2011], Perl [Moertel, 2011], Ruby [Yeh, 2011] and JavaScript [Thompson, 2011] among others. These tools are all about test data generation. By generating high amounts of data and checking it against a given property, it is expected to cover a wide range of the input domain, thus increasing the changes of finding more faults [Hughes, 2009]. While there is no validation for the Java implementation of QUICKCHECK, there is a report on the effectiveness of QUICKCHECK for Haskell [Claessen and Hughes, 2000].

QUICKCHECK for Java [Open Source Contributors, 2011a] is one implementation of QUICKCHECK for the Java language that allows for the creation of PUT cases. Recall that PUT is to run the same test case with different input values. This implementation includes several statistical distributions allowing users to better guide data generation.

If applied like a third party framework, QUICKCHECK requires test cases to explicitly call its features. In some occasions it might be necessary to implement a custom data generator, which QUICKCHECK allows and provides an extensive library for. Listing 3.1 illustrates an example implementation of a generator. In this implementation a Stack instance is created with a random size between 1 and 100, and then filled with random integers.

¹This requires a different version of the tool.

Language	Java and .NET
Price	Free under the MIT licence (see MIT [2011])
Level of Tests	Unit level
Scope	Test case generator
Adequacy Criterion	Fault detection
Incremental	Not influenced by existing test cases
Previously Generated Tests	Overwrites or complements previously generated test cases
Integration With Manual Tests	Possible to co-exist
	Comes as an Eclipse IDE plugin and a command line
Integration With Development Tools	tool. There is however, no integration with Apache Ant, Apache Maven and Netbeans
Integration With Development Tools Code or Binary Level	tool. There is however, no integration with Apache Ant, Apache Maven and Netbeans Works at the binary level
Integration With Development Tools Code or Binary Level Configuration Effort	tool. There is however, no integration with Apache Ant, Apache Maven and Netbeans Works at the binary level Little effort in using the Eclipse plugin. The standalone command line tool requires a proper setup of the class path variable

Table 3.2: Summary for RANDOOP.

```
public class StackGenerator implements Generator<FastStack<Integer>> {
    public FastStack<Integer> next() {
        Random rand = new Random();
        int size = 1 + rand.nextInt(100);
        FastStack<Integer> fs = FastStack.newInstance();
        for (int i = 0; i < size; i++) {
            fs.push(rand.nextInt());
        }
        return fs;
    }
}</pre>
```

Listing 3.1: Implementation of a random Stack generator

Implemented generators can then be applied within a test case as illustrated in Listing 3.2. In this test case, the generated Stack instance is tested through some common operations and checked for conformity. Then the Stack instance is regenerated and the process starts all over again.

```
@Test
public void testNormal() {
  FastStack<Integer> stack;
  for (int i = 0; i < NUMBER_OF_RUNS; i++) {</pre>
    stack = sg.next();
    int currentTop = stack.peek();
    int currentSize = stack.size();
    final int num = 100;
    for (int j = 0; j < num; j++) {</pre>
        stack.push(j);
      assertEquals(j, stack.peek());
    }
    assertEquals(currentSize + num, stack.size());
    assertEquals(1, stack.peek(currentSize + 1).intValue());
    for (int j = 0; j < num; j++) {</pre>
        stack.pop();
    }
    assertEquals(currentTop, stack.peek().intValue());
    assertEquals(currentSize, stack.size());
  }
}
```

Listing 3.2: Application of the Stack generator in a test case

ЈСнеск

JCHECK [Open Source Contributors, 2011b] is another Java implementation of the previously mentioned QUICKCHECK combinator library [Claessen and Hughes, 2000]. This implementation however presents a different approach to how the test cases are written. Instead of adding a loop to iterate over the randomly generated values, JCHECK features its own test runner, which is compatible with the JUnit framework. This allows test cases to receive input parameters and run the test case with the received parameters. This implementation of QUICKCHECK lacks many of the data generation features present in QUICKCHECK for Java.

Much like QUICKCHECK for Java, to apply JCHECK the test cases must include its functionalities. However, there is little difference while implementing a data generator, the difference being that the random value generator and the scope size are already passed on to the generator as parameters. A PUT written with JCHECK is illustrated in Listing 3.3. In this case, the generator is specified through an annotation provided by JCHECK. Notice the lack of a loop statement to refresh the instance under test, this being accomplished by the runner that comes packaged with JCHECK.

Language	Java
Price	Free under the Apache 2.0 licence (see ASF [2011])
Level of Tests	Unit level
Scope	Parameterised test case framework
Adequacy Criterion	Input domain coverage
Incremental	Not influenced by existing test cases
Previously Generated Tests	Impossible to replay previous instantiated tests
Integration With Manual Tests	Possible to co-exist
Integration With Development Tools	Only available as a third party library
Code or Binary Level	Works at the source level
Configuration Effort	To use the library it is only required to add it to the build path and use its features
Quality of Produced Tests	Does not produce the code for the test case, it only in- stantiates a test case with different input values

Table 3.3: Summary for QUICKCHECK.

```
@Test
@Generator(klass = FastStack.class, generator = StackGenerator.class)
public void testNormal(FastStack<Integer> stack) {
  int currentTop = stack.peek();
  int currentSize = stack.size();
  final int num = 100;
  for (int i = 0; i < num; i++) {</pre>
      stack.push(i);
  }
  assertEquals(currentSize + num, stack.size());
  assertEquals(num - 1, stack.peek().intValue());
  for (int i = 0; i < num; i++) {</pre>
   stack.pop();
  }
  assertEquals(currentTop, stack.peek().intValue());
  assertEquals(currentSize, stack.size());
}
```

Listing 3.3: Application of the Stack generator in a parameterised test case

Language	Java
Price	Free under the CPL licence (see IBM [2011])
Level of Tests	Unit level
Scope	Parameterised test case framework
Adequacy Criterion	Input domain coverage
Incremental	Not influenced by existing test cases
Previously Generated Tests	Impossible to instantiate previous test cases
Integration With Manual Tests	Possible to co-exist
Integration With Development Tools	Only available as a third party library
Code or Binary Level	Works at the source level
Configuration Effort	To use the library it is only required to add it to the build path and use its features
Quality of Produced Tests	Does not produce the code for the test case, it only in- stantiates a test case with different input values

Table 3.4: Summary for JCHECK.

FEED4JUNIT

FEED4JUNIT is another PUT generation framework. It is developed and maintained by Volker Bergmann [Volker Bergmann, 2011]. This tool allows users to specify an external *.csv file to feed input data and expect values to a given test case. This is especially useful since anyone from software analysts to developers can build a complete *.csv file with all the important input values to be tested. It also allows developers to produce equivalence class partition test cases, this meaning that it is possible to define strict input sets for each test case. FEED4JUNIT is capable of exploring those input sets by following a combinatorial strategy. Consider the example in Listing 3.4, in which two sets for each input parameter are defined, FEED4JUNIT would combine the sets and test them thoroughly.

There is little difference between tests written with this framework and the ones written with the previously described tools, QUICKCHECK for Java and JCHECK. The differences in implementing custom generators and writing test cases are mostly aesthetic.

T2 Framework

T2 is a TCG framework developed by Wishnu Prasetya and collaborators [Prasetya et al., 2008]. This tool is a trace-based random test generation tool, similar to RANDOOP. However, there is a difference: T2 allows users to provide in-code specifications through the assert keyword of the Java programming language. Furthermore, this tool integrates with JUnit. Downside of this tool, is the fact that it can not test non-deterministic behaviour. For instance, classes that make use of randomly generated values or multi-threaded behaviour can not be tested with T2. A study of the effectiveness of T2 [Prasetya, 2011], reported that T2 can reduce the cost of testing software. Its speed and achieved coverage are its stronger points.

```
@RunWith(Feeder.class)
public class AddTest {
  0Test
  @Unique
  public void testAdd(
      @Values("-2147483648,-6,-1,0,1,42,2147483647") int param1,
      @Values("-2147483648,-6,-1,0,1,42,2147483647") int param2) {
    try {
      int result = MyUtil.add(param1, param2);
    } catch (Exception e) {
      \ensuremath{{\prime\prime}}\xspace application exceptions, fail on runtime exceptions
      // like NullPointerException
      if (e instanceof RuntimeException)
        throw e;
    }
  }
}
```

Listing 3.4: Example of a Feed4JUnit equivalence partition test.

Language	Java
Price	Free under GPLv2 (see GNU [2011]) but depends in other components that are not entirely GPL (see Volker Bergmann 2011 for more details)
Level of Tests	Unit level
Scope	Parameterised test case framework
Adequacy Criterion	Input domain coverage
Incremental	Not influenced by existing test cases
Previously Generated Tests	Impossible to instantiate previous test cases, however it is possible to check which inputs were passed on to the test case
Integration With Manual Tests	Possible to co-exist
Integration With Development Tools	Only available as a third party library
Code or Binary Level	Works at the source level
Configuration Effort	To use the library it is only required to add it to the build path and use its features
Quality of Produced Tests	Does not produce the code for the test case, it only in- stantiates a test case with different input values

Table 3.5: Summary for FEED4JUNIT.

To apply T2, it is recommended that the source code should contain some specifications, otherwise the test process will only check for run-time exceptions. The specifications should be introduced through the assert keyword, either in the desired method or in a separate method with the same name appended with a _spec suffix, as illustrated in Listing 3.5. To invoke the T2 runner there are two possible ways. One way is to invoke the runner through the command line. However, it is also possible, and recommended, to invoke T2 within a JUnit test case. Running T2 from a test case allows other tools to measure the structural coverage that T2 managed to achieve. Listing 3.6 illustrates how T2 can be invoked from JUnit.

```
public class FastStack<E> implements Iterable<E> {
    private final List<E> mEntries = Lists.newArrayList();
    ...
    public void push(E aElement) {
        mEntries.add(aElement);
    }
    public void push_spec(E newElem) {
        assert newElem != null : ``PRE'';
        int size = this.size();
        this.push(newElem);
        assert (!this.isEmpty()) : ``POST'';
        assert (this.contains(newElem)) : ``POST'';
        assert (this.size() == size + 1) : ``POST'';
    }
    ...
}
```



```
@Test
public void testClass() {
   Sequenic.T2.Main.Junit(FastStack.class.getName()
   + ``--nmax=2000 --lenexec=5 --violmax=8 --elemty=java.lang.Integer'');
}
```

Listing 3.6: Example of a test case invoking T2.

JMLUNITNG

JMLUNITNG [Zimmerman and Nagmoti, 2011; Zimmerman, 2011], is a TCG tool for JML-annotated source code. The tool is developed and maintained by Daniel Zimmerman. It aims to support the most recent features of the Java language, like generics and enumerated types. It leverages JML specifications to generate test oracles, it also attempts to use reflection mechanisms to generate test data, and also provides users the possibility to introduce their own input values. Currently JMLUNITNG relies on the TestNG framework for running the generated test cases and a compiler of the JML language (either jmlc or jml4c). Additionally, JMLUNITNG also depends on the next major version of the JDK, OpenJDK 7, which is, at the time of writing, under development.

Language	Java
Price	Free under GPLv3 licence (see GNU [2011])
Level of Tests	Unit level
Scope	Test case generator
Adequacy Criterion	Specification validation / input domain coverage
Incremental	Not influenced by existing test cases
Previously Generated Tests	Previously generated tests can be re-run, with the aim of fixing bugs
Integration With Manual Tests	Possible to co-exist
Integration With Development Tools	Only available as a third party library
Code or Binary Level	Works at the source level
Configuration Effort	To use the library it is only required to add it to the build path and use its features. Running it from the command line the class path needs to be properly configured. The specifications need to be added by other means
Quality of Produced Tests	Does not produce the code for the test cases, it only in- stantiates a test case with different input values

Table 3.6: Summary for T2.

The first step in applying this tool is to properly annotate the source code with JML specifications, whereafter the classes need to be compiled with a special compiler: jmlc if the source code does not require newer Java features or with jml4c if it requires newer features, like generic types. Then JMLUNITNG needs to run on the classes under test with java -jar jmlunitng.jar path-list command. This command will generate the required java classes which need to be compiled and executed with OpenJDK 7. Before compiling the generated classes, they can be edited to provide input values.

Language	JML annotated Java
Price	Free, not open source
Level of Tests	Unit level
Scope	Test case generator
Adequacy Criterion	Specification validation / input domain coverage
Incremental	Not influenced by existing test cases
Previously Generated Tests	Only generates test oracle, test cases can be added by extending the input set
Integration With Manual Tests	Possible to co-exist
Integration With Development Tools	Only available as a third party tool, there is no integration with any other development tool
Code or Binary Level	Works at the source level
Configuration Effort	This tool requires OpenJDK 7 installed and set up as the main JVM. The JML compilers, $jmlc$ and $jml4c$, need to be set up, and lastly the class path variable needs to be set up to include JMLuniNG
Quality of Produced Tests	Does not produce the code for the test cases, it only in- stantiates a test case with different input values

Table 3.7: Summary for JMLUNITNG.

JAVA PATHFINDER

JAVA PATHFINDER is a tool developed by researchers at the NASA Ames Research Centre [Visser and Mehlitz, 2005; NASA Ames Research Center, 2011]. This tool is more a model checker for Java bytecode than a real TCG. This however does not mean that it can not be used to test software. In fact the type of model checking that this tool performs is commonly mistaken with systematic testing. The truth is that JAVA PATHFINDER does a little more than traditional TCG tools. Tests generated with traditional tools only cover one execution path at a time, and some times some faults may go on undetected, while on the other hand model checking systematically checks all paths. This increases the chances of finding more faults (see Figure 3.2).

JAVA PATHFINDER is a virtual machine on top of the regular JVM. This new virtual machine adds some model checking traits like backtracking and state matching among others to the original JVM.



Figure 3.2: Software Testing vs. Model Checking [NASA Ames Research Center, 2011]

JAVA PATHFINDER was successfully applied to check random algorithms [Zhang and van Breugel, 2010] and network protocols [Martínez and Jiménez, 2008] with success. There are, however, some limitations, for example it is only meant for non-deterministic software as quoted from the project page [NASA Ames Research Center, 2011]:

A word of caution: if you have a strictly sequential program with only a few well defined input values, you are probably better off writing a few tests - using JPF won't tell you much.

PATHFINDER is also unable to check Java Native Interface methods, this is, JAVA PATHFINDER can not test most I/O operations that the language supports and in which most software currently build on. Out-of-the-box, JAVA PATHFINDER checks the following properties, leaving the possibility to implement new properties:

- **Deadlocks** gov.nasa.jpf.jvm.NotDeadlockedProperty for every non-end state, test if there is any runnable thread left;
- **Assertion Violation** gov.nasa.jpf.jvm.NoAssertionViolatedProperty test if any assertion expression has been violated;
- **Uncaught Exceptions** gov.nasa.jpf.jvm.NoUncaughtExceptionsProperty test if any exception was not handled inside the software.

To apply JAVA PATHFINDER to a software under test, it first needs to to be configured to define the intended software properties. Listing 3.7 illustrates a sample configuration, the target property defining the class where the software entry point is located, DiningPhil in the example. The search .class property defines the search strategy to be applied in the process, a breath-first approach in the example. Listing 3.8 illustrates the class that will be targeted. This class consists of the

traditional dinning philosophers problem, in which there are five philosophers and five forks. A philosopher is either thinking or eating, when a philosopher wants to eat it needs to take both the fork to its left and to its right. To start the JAVA PATHFINDER process the user just needs to invoke jpf <path-to-jpf-configuration-file> from the command line.

target = DiningPhil

```
search.class = .search.heuristic.BFSHeuristic
```

Listing 3.7: Example JPF configuration.

Language	Java
Price	Free under the NOSA licence (see NASA [2011])
Level of Tests	System level
Scope	Java bytecode model checker
Adequacy Criterion	Property validation / structural coverage
Incremental	Not influenced by existing test cases
Previously Generated Tests	Does note generate test cases
Integration With Manual Tests	Possible to co-exist
Integration With Development Tools	Plugins available for Eclipse and Netbeans, there is cur- rently no support for Ant or Maven
Code or Binary Level	Works at the bytecode level
Configuration Effort	Configuration depends on each system system under test
Quality of Produced Tests	Does not produce test cases

Table 3.8: Summary for JAVA PATHFINDER.

UNITCHECK

UNITCHECK is a plugin a tool that was developed by Michal Kebrt [Kebrt and Sery, 2009]. This tool means to apply model checking techniques, through JAVA PATHFINDER to unit test cases, aiming to bring high quality model checking to the general software development process. UNITCHECK recognises JUnit tests and exhaustively explores their reachable state space including all admissible thread interleaving. This tool was used to check the behaviour of two test classes that lean on threaded behaviour to reveal faults [Kebrt and Sery, 2009].

Udita

UDITA [Gligoric et al., 2010] is an extension to the standard Java language, that allows its users to describe their test cases. Relying in a modified version of Java PATHFINDER tool. This tool allows

```
import gov.nasa.jpf.jvm.Verify;
public class DiningPhil {
  static class Fork {}
  static class Philosopher extends Thread {
   Fork left;
    Fork right;
    public Philosopher(Fork left, Fork right) {
     this.left = left;
      this.right = right;
     start();
    }
    public void run() {
      // think!
      synchronized (left) {
       synchronized (right) {
         // eat!
        }
  static final int N = 6;
  public static void main(String[] args) {
   Verify.beginAtomic();
   Fork[] forks = new Fork[N];
    for (int i = 0; i < N; i++)</pre>
     forks[i] = new Fork();
    for (int i = 0; i < N; i++)</pre>
     new Philosopher(forks[i], forks[(i + 1) % N]);
    Verify.endAtomic();
  }
}
```

Listing 3.8: The classic dining philosophers problem.

Language	Java
Price	Free, no licensing information is available
Level of Tests	Unit level
Scope	Unit level model checker
Adequacy Criterion	Property validation / structural coverage
Incremental	Not influenced by previous executions
Previously Generated Tests	Does note generate test cases
Integration With Manual Tests	Required
Integration With Development Tools	Available as an Eclipse plugin and as an Ant Task, cur- rently there is no support for Netbeans IDE or Maven
Code or Binary Level	Works at the binary level
Configuration Effort	Little effort in configuring the Eclipse plugin, running the Ant task however requires some understanding of Ant and how to set up the class path variable
Quality of Produced Tests	Does not produce test cases

Table 3.9: Summary for UNITCHECK.

users to describe either test oracles or the input domain for the test cases. The tool will generate input values or appropriate test cases to verify the correction of the oracles. These inputs could be used at either system testing level or at the unit testing level. It has been evaluated on the JAVA PATHFINDER and Eclipse IDE to reveal some flaws in both projects [Gligoric et al., 2010].

ECLAT

ECLAT is very similar to RANDOOP and has been developed by Carlos Pacheco and Michael Ernst [Pacheco and Ernst, 2005]. The major difference is that ECLAT integrates with DAIKON in order to produce an operational model of the system under test. This model is then used to generate test cases, by filtering out redundant test cases and improving the ability to check for faults. The developer can then check the failing test cases and decide if the test case is really revealing a fault of the software or not.

The tutorial available at [Pacheco and Ernst, 2005] provides an easy to understand example of usage. The basic usage is to tell EcLAT which classes should it generate test cases for and an example usage of those classes, for example existing test cases over those classes. A typical command line would be java eclat.textui.Main generate-inputs --create-regression-suite --test ubs/BoundedStack.java ubs.BoundedStackJunitTest which will generate a test suite. The resulting test suite can be found in a folder named eclat-src.

Рех

PEx is a tool developed by Nikolai Tillmann and Peli de Halleux and other staff at Microsoft Research labs [Tillmann and de Halleux, 2008]. PEx leverages concolic execution to execute pro-

Language	Java
Price	Free under the NCSA licence (see University of Illinois [2011])
Level of Tests	System / unit level
Scope	Input generator / Oracle generator
Adequacy Criterion	Input domain coverage
Incremental	Not influenced by previous executions
Previously Generated Tests	Not applicable
Integration With Manual Tests	Possible to coexist
Integration With Development Tools	Does not integrate with any software development tool
Code or Binary Level	Works at the binary level
Configuration Effort	Impossible to say, the instructions provided did not lead to a successful execution of the tool as described
Quality of Produced Tests	Does not produce test cases

Table 3.10: Summary for UDITA.

Java
Free, in spite of, the source code not being currently available
Unit level
Test case generator
Fault detection
Existing test cases improve the result
Can be ran to check remaining faults
Possible to co-exist
Only available as a third party library
Works at the source level
To use the library it is only required to add it to the build path and use its features
Produces complete test cases, with inputs and test or- acles.

Table 3.11: Summary ECLAT.

grams and capture their behaviour, generating a PUT and passing test inputs on to fully exercise the function under test and detect errors.

This tool is fully integrated with the Visual Studio IDE from Microsoft. After installing PEx the only thing to do is to press a single button and PEx will automatically start generating test cases.

Language	.NET
Price	Free for Visual Studio
Level of Tests	Unit level
Scope	Test case generator
Adequacy Criterion	Structural coverage / Specification validation (if com- bined with Code Contracts)
Incremental	Not influenced by previously generated tests
Previously Generated Tests	May remain for future notice or can be overwritten
Integration With Manual Tests	Possible to co-exist
Integration With Development Tools	Integrates with Microsoft Visual Studio, only
Code or Binary Level	Works at the binary level
Configuration Effort	Only installation is required
Quality of Produced Tests	Produces parameterised tests cases with test oracles and appropriate data generators

Table 3.12: Summary for PEX.

3.2.2 Interesting Tools for Software Testing

This section describes interesting tools in the software testing context, which are not TCG tools.

Korat

KORAT is a tool that automatically generates structurally complex test inputs for the Java programming language. It is being developed by researchers at MIT and other universities [Milicevic et al., 2007]. Users must provide a predicate that describes a proper instantiation of the desired structure and a finitisation criterion to limit the scope. However, KORAT does not integrate with other testing frameworks like JUnit *out-of-the-box*.

To use KORAT, users must first provide at least two important methods to the target structure. These methods consist of (*a*) an imperative predicate that specifies the desired structural constraints; and (*b*) a finitisation criterion that bounds the desired test input size. An imperative predicate is a method that checks that all the desired properties about the structure are kept. It does not take any parameters and should always return a Boolean value, true in case the structure is valid or false otherwise. A finitisation is a method that tells the tool when to stop generating input values. In a way it bounds the input domain to prevent from memory exhaustion.

Consider the implementation of a binary tree, illustrated in Listing 3.9. This structure holds up to two nodes and each node can hold up to two other nodes and so on. A possible predicate over this structure would be Listing 3.10, that prevents the occurrence of cycles in the structure. The finitisation method (Listing 3.11), as previously mentioned, simply defines the size of the structure to generate. After enriching the desired structure with both these methods, KORAT can be applied. Currently, it is only available as a command line tool. To invoke KORAT on the presented examples, the following command could be used: java korat.Korat --visualize --class korat.examples .binarytree.BinaryTree --args 5, 5, 5. This command would produce some visualisations of the intended structure.

```
public class BinaryTree {
    public static class Node {
        Node left;
        Node right;
    }
    private Node root;
    private int size;
}
```



```
public boolean repOK() {
   if (root == null)
       return size == 0;
    // checks that tree has no cycle
   Set visited = new HashSet();
   visited.add(root);
   LinkedList workList = new LinkedList();
   workList.add(root);
   while (!workList.isEmpty()) {
       Node current = (Node) workList.removeFirst();
        if (current.left != null) {
           if (!visited.add(current.left))
               return false;
           workList.add(current.left);
        if (current.right != null) {
            if (!visited.add(current.right))
                return false;
            workList.add(current.right);
        }
    }
    // checks that size is consistent
   return (visited.size() == size);
}
```

Listing 3.10: Example binary tree imperative predicate.

So by applying KORAT to the previously illustrated binary tree it is possible to quickly generate various instances of the tree, one of them being illustrated in Figure 3.3.

```
public static IFinitization finBinaryTree(int nodesNum, int minSize, int maxSize) {
    IFinitization f = FinitizationFactory.create(BinaryTree.class);
    IObjSet nodes = f.createObjSet(Node.class, nodesNum, true);
    f.set("root", nodes);
    f.set("Node.left", nodes);
    f.set("Node.right", nodes);
    IIntSet sizes = f.createIntSet(minSize, maxSize);
    f.set("size", sizes);
    return f;
}
```





Figure 3.3: Example visualisation of a BinaryTree of size 5.

Language	Java
Price	Free under the GPLv2 licence (see GNU [2011])
Level of Tests	Not applicable
Scope	Input generator
Adequacy Criterion	Input domain coverage
Incremental	Not influenced by previously generated inputs
Previously Generated Tests	Does note generate test cases
Integration With Manual Tests	No integration
Integration With Development Tools	Does not integrate with IDEs nor other development tool
Code or Binary Level	Works at the bytecode level
Configuration Effort	KORAT only needs to be added to the class path; the im- perative predicates and finitisation methods need to be written
Quality of Produced Tests	Does not produce test cases

Table 3.13: Summary for KORAT.

REASSERT

REASSERT is helpful as a maintenance tool. It was developed by Brett Daniel and Darko Marinov, [Daniel et al., 2011]. This tool helps its users in correcting broken tests. It must be used with caution as the tool will not know if developers have introduced a new bug into the function or just altered its behaviour to include new functionality. This tool executes the function whose tests got broken with the last revision and examines its new behaviour, then it suggests some corrections to the failing test that make it pass again, be it by changing an assertion or an input value. Figure 3.4 illustrates the typical usage of REASSERT

DAIKON

DAIKON is an invariant detection tool developed at MIT by Michael Ernst and collaborators [Ernst et al., 2007]. This tool is based in dynamic analysis of applications, this is, it detects likely program invariants based on values observed from actual executions of the software. Not directly related to TCG, DAIKON can be an intermediate step of it, since invariants may be used to aid the testing process, either manually or automatically. Manually, in the sense that it may help developers to better understand how a given function or method should behave. TCG tools could leverage the invariants to better guide their generation process.

To apply DAIKON, users should provide some executions over the intended software, which can be regular main methods exercising certain parts of the software, or regular JUnit tests. At the moment DAIKON is only available as a command line tool. To invoke it on an existing test suite the command java daikon.Chicory --daikon TARGET_CLASS should be enough. Then the tool will execute the target classes and capture the execution traces that are then used to infer the likely invariants. DAIKON is also able to annotate source code files with the detected invariants with the



- Confirm	proposed fixes		
≽ Structure Compare			
≂ 🗁 cart	▲ ₩		
🗎 🛛 CartTest.java.fix / CartTest.java			
📄 Text Compare	😥 😓 📣 🕸 🖓		
Original Test File	Fixed Test File		
public void testRedPenCoupon() {	public void testRedPenCoupon() {		
Cart cart = new Cart();	Cart cart = new Cart();		
cart.addProduct(new RedPen());	<pre>cart.addProduct(new RedPen());</pre>		
<pre>cart.addProduct(new RedPen()); cart.addProduct(new RedPen());</pre>			
<pre>cart.addCoupon(new AnniversaryCoupon()); cart.addCoupon(new AnniversaryCoupon());</pre>			
assertEquals (3.0, cart.getTotalPrice()); assertEquals (6.0, cart.getTotalPrice());			
assertEquals (assertEquals (
"Discount: -\$3.00, Total: \$3.00",	"Discount: \$0.00, Total: \$6.00",		
cart.getPrintedBill());	<pre>cart.getPrintedBill());</pre>		
}	}		
• *****			
0	C <u>o</u> nfirm Cancel		

(b) Confirm suggested corrections.

Figure 3.4: Applying REASSERT

Language	Java
Price	Free under the NCSA licence (see University of Illinois [2011])
Level of Tests	Unit level
Scope	Unit test correction tool
Adequacy Criterion	Not applicable
Incremental	Not applicable
Previously Generated Tests	Not applicable
Integration With Manual Tests	Possible, since it repairs test cases
Integration With Development Tools	Plugins available for Eclipse, only
Code or Binary Level	Works at the bytecode level
Configuration Effort	Only the plugin installation is required
Quality of Produced Tests	Does not produce test cases

Table 3.14: Summary for REASSERT.

following command java daikon.tools.jtb.Annotate INVARIANT_FILE TARGET_SOURCE_FILE.

AUTOMATIC JUNIT CREATION TOOL

AUTOMATIC JUNIT CREATION TOOL [Whitney, 2011] is a tool that allows the easy creation of JUnit test cases. This tool allows users to interact more closely with the methods under test by allowing the user to specify which inputs to pass to the method under test and what behaviour to expect. To put it simply, AUTOMATIC JUNIT CREATION TOOL is a JUnit test stub generator with the ability to interact directly with the JUnit test being generated. Once all the methods are properly set up the tool allows the user to export a complete JUnit test file. Perhaps the most interesting feature provided by the tools is the ability to automatically identify branches in the control flow of a method and allow the user to specify independent test cases for them.

This tool is only available as a standalone GUI application. In order to use the tool a user should first load a java class file using the menu-bar (File \rightarrow New Test) and choose the file for which to stub out new test cases. The buttons at the bottom of the window (see Figure 3.5) can be used to cycle through the available methods and gradually build each test case, passing on inputs and defining appropriate test assertions. When all the methods have proper inputs and assertions the user may export the test cases to a Java file to be used by external tools like JUnit.

JWalk

JWALK is a tool that allows lazy systematic unit testing of Java classes. It was developed by Anthony Simons [Simons, 2007] and is currently licensed by the University of Sheffield. This tool preforms bounded exhaustive testing of any compiled java class and it tests it for conformance to a lazy specification inferred by the tool at run-time from the code.

Language	Various
Price	Free under the MIT licence (see MIT [2011])
Level of Tests	Not applicable
Scope	Invariant Detector
Adequacy Criterion	Not applicable
Incremental	The more executions available, the more trustworthy the invariants are
Previously Generated Tests	Not applicable
Integration With Manual Tests	Possible to leverage unit tests as executions
Integration With Development Tools	Does not integrate with any IDE or software developing tool
Code or Binary Level	Works at the binary level
Configuration Effort	Requires proper set up of the class path variable, where- after it is fairly simple to use
Quality of Produced Tests	Does not produce test cases

Table 3.15: Summary for DAIKON.

Figure 3.5: Automatic JUNIT CREATION TOOL main window.

Language	Java
Price	Free under the GPLv3 licence (see GNU [2011])
Level of Tests	Unit level
Scope	Test case stub generator
Adequacy Criterion	Branch coverage
Incremental	Not influenced by existing test cases
Previously Generated Tests	Can not be opened by the tool again
Integration With Manual Tests	This tool facilitates writing of test cases
Integration With Development Tools	Only available as a standalone application
Code or Binary Level	Works at the source level
Configuration Effort	No configuration effort
Quality of Produced Tests	Produces complete test cases, with inputs and asser- tions

Table 3.16: Summary for Automatic JUNIT CREATION TOOL.

JWALK is available as a Java GUI application where the user is allowed to set the parameters for the test process. One of three strategies can be chosen: protocol, algebra and states. And one of three modalities: inspect, explore and validate and lastly the depth of the test process. The protocol strategy creates test sequences corresponding to every possible interleaved ordering of the public constructors and methods in the class under test. The algebra strategy creates sequences that drive the test object into different concrete states. The last strategy, states, creates sequences to cover most machine states of the class.

The inspect modality preforms a static analysis of the class to reveal the public API of the the test class or the state space of the class. The explore modality preforms a similar analysis, but dynamically instead. The validate modality creates and executes test sequences and validates the outcome of the sequence at hand. The test depth defines the bound the algorithms should use for bounded exhaustive testing. It turns out that this tool does not export any usable JUnit files that can be used by other tools.

3.3 A Case Study - Checkstyle

This section starts by presenting, Checkstyle, a piece of software chosen as target for the application of the testing tools described thus far. In this way we intend to "test such testing tools" checking how they behave and drawing conclusions on their usefulness.

3.3.1 Subject of Study: Checkstyle

To preform our study, Checkstyle was chosen as the target of the tools under review. Checkstyle was developed by Oliver Burn [Burn, 2011] and other freelance contributors. The purpose of this tool is to help developers in writing Java code that adheres to a given coding standard. It can

Language	Java
Price	Free for academic uses, not open source
Level of Tests	Unit level
Scope	Test case regeneration tool
Adequacy Criterion	Bounded exhaustive testing
Incremental	Previously generated test oracles can be used in gener- ation of new tests
Previously Generated Tests	Stored as a binary file
Integration With Manual Tests	Possible to co-exist
Integration With Development Tools	Only available as a GUI application
Code or Binary Level	Works at the binary level
Configuration Effort	JWALK requires a licence under which the jar file executes
Quality of Produced Tests	Not producing the code for the test cases, it only pro- duces a binary file that the tool uses to regenerate tests

Table 3.17: Summary for JWalk.

be configured to check virtually every coding style, but it defaults to the standard Java coding conventions [Sun Microsystems, Inc., 1999] that comes packaged with the tool. Other reasons for the success of Checkstyle, are easy the integration with other Java development tools like the Eclipse IDE and the Apache Maven project manager. It is important to note that Checkstyle performs several I/O operations, in reading: *a*) Checkstyle configuration files; *b*) Desired coding styles; and *c*) source files to check. There is also the printing of a report that can be analysed by the users of the tool, and information about the process that is sent to the default system logger.

This application was chosen for our case study because it is a common Java application, and for its history in software engineering research [Cornelissen and Moonen, 2007; Zaidman et al., 2007, 2008; Lubsen et al., 2009]. Before digging in to the case study itself, some information about the application was collected, especially information about its source code. First, an overview of Checkstyle has shown that the tool is composed of 22 packages and over 300 classes (Table 3.18). The first line of Table 3.18 presents the total metrics for the project while the other lines provide an average value for each level. There is also an existing test suite for Checkstyle, which allows some comparisons to be made.

Due to time limitations, it was decided to conduct this study in a smaller, but also relevant, part of the tool. In order to locate which part of the application should be used while evaluating TCG tools, a Fan-In/Fan-Out table of the packages of Checkstyle was derived (Table 3.20). The most relevant package was found to be com.puppycrawl.tools.checkstyle.api, due to the number of packages that depend on it. The existing test cases for the package cover about 3.5 % of the whole application as detailed in Table 3.19.

	Lines of Javadoc	Non Comment Source Statements	Methods	Classes	Packages
Project	1,674	20,781	2,240	327	22
Package	76	945	102	15	
Class	5	64	7		
Method	1	9			

Table 3.18: Information collected from Checkstyle.

Element	Coverage (%)	Covered Instructions	Missed Instructions	Total Instructions
Target Package	73.2	3,896	1,430	5,326
src/checkstyle	81.3	36,791	8,480	45,271
target/generated-sources/antlr	83.5	20,348	4,019	24,367
Totals	31.4	57,139	12,499	69,638
Average	43.4	31.4%	68.6%	100.00%

Table 3.19: Coverage results achieved by the manual test cases.

3.3.2 Tools Tested

The TCG tools that were successfully applied to the subject of our study, <code>com.puppycrawl.tools</code>.checkstyle.api, were: RANDOOP, QUICKCHECK for Java, JCHECK, FEED4JUNIT, T2, DAIKON, AUTOMATIC JUNIT CREATION TOOL, REASSERT and JWALK. In this section we will discuss the results obtained from exercising such tools.

Starting with RANDOOP, it was decided to run in two time spans, 2.5 and 5 minutes. Table 3.21 shows the coverage results obtained with both the test suites. Clearly there is little gain in increasing the running time from 2.5 to 5 minutes. Even with 5 minutes RANDOOP was unable to achieve a coverage result nearly as remarkable as the manually written test cases.

Let us now proceed by discussing the three PUT tools that were covered by this study, QUICKCHECK for Java, JCHECK and FEED4JUNIT. These tools were applied successfully to the target package and all of them achieved the same line coverage as the manually written tests. They have, however, allowed the coverage over the input domain to increase, by instantiating the test cases for several inputs.

The T2 Framework was also applied with success to a class of the target package, com.puppycrawl .tools.checkstyle.api. Because this tool requires some specifications in order to reveal its full potential, it was decided to produce specifications for one class (FastStack). This tool was able to cover 12.8% of the total instructions in comparison with the 15.9% achieved by the manually written test cases. This is not a bad result considering it was done automatically with just some simple specifications.

The application of DAIKON used only a fraction of the test suite of the case study. It was possible to infer some invariants from test cases, an example of which can be found in Listing 3.12. An

Package	Fan-In	Fan-Out
com/puppycrawl/tools/checkstyle	96	166
com/puppycrawl/tools/checkstyle/api	993	351
com/puppycrawl/tools/checkstyle/checks	180	211
com/puppycrawl/tools/checkstyle/checks/annotation	20	49
com/puppycrawl/tools/checkstyle/checks/blocks	17	31
com/puppycrawl/tools/checkstyle/checks/coding	82	232
com/puppycrawl/tools/checkstyle/checks/design	11	38
com/puppycrawl/tools/checkstyle/checks/duplicates	31	56
com/puppycrawl/tools/checkstyle/checks/header	5	10
com/puppycrawl/tools/checkstyle/checks/imports	29	58
com/puppycrawl/tools/checkstyle/checks/indentation	118	145
com/puppycrawl/tools/checkstyle/checks/javadoc	22	77
com/puppycrawl/tools/checkstyle/checks/metrics	26	55
com/puppycrawl/tools/checkstyle/checks/modifier	2	7
com/puppycrawl/tools/checkstyle/checks/naming	27	48
com/puppycrawl/tools/checkstyle/checks/regexp	24	34
com/puppycrawl/tools/checkstyle/checks/sizes	14	41
com/puppycrawl/tools/checkstyle/checks/whitespace	26	54
com/puppycrawl/tools/checkstyle/doclets	10	10
com/puppycrawl/tools/checkstyle/filters	24	66
com/puppycrawl/tools/checkstyle/grammars	10	8
com/puppycrawl/tools/checkstyle/gui	153	173
Total	1,920	1,920

Table 3.20: Fan-In and Fan-Out by package of Checkstyle.

Element	Test Suite	Coverage (%)	Covered Instructions	Missed Instructions	Total Instructions
Taraat	Randoop 2.5 min	38.60	2,057	3,269	
Deckage	Randoop 5 min	39.30	2,092	3,234	5,326
гаскауе	Manual Tests	73.20	3,896	1,430	

Table 3.21: RANDOOP generated tests for the target package compared with manually written tests.

example of an annotated method with the extracted invariants is illustrated in Listing 3.13.

```
com.puppycrawl.tools.checkstyle.api.FastStack.clear():::ENTER
Variables: this this.mEntries this.mEntries.getClass()
this != null
this.mEntries != null
com.puppycrawl.tools.checkstyle.api.FastStack.clear():::EXIT
Variables: this this.mEntries this.mEntries.getClass() orig(this) orig(this.
mEntries) orig(this.mEntries.getClass())
this.mEntries == \old(this.mEntries)
this.mEntries != null
this.mEntries != null
```

Listing 3.12: Two Invariants inferred by DAIKON.

```
public class FastStack<E> implements Iterable<E> {
    ...
    /*@ requires this != null; */
    /*@ requires this.mEntries != null; */
    /*@ ensures this.mEntries != null; */
    /*@ ensures \typeof(this.mEntries) == \old(\typeof(this.mEntries)); */
    public void clear() {
        mEntries.clear();
    }
    ...
}
```



Albeit useful, the invariants detected by DAIKON are incomplete in general. Take for instance the annotated method in Listing 3.14, which only take one of three possible values: 0, 1 or 100. This is obviously incomplete, since there is nothing preventing a Stack instance from having any particular size.

AUTOMATIC JUNIT CREATION TOOL was also successfully applied to the case study. The work flow of this tool is not much of an improvement over manually writing the test cases. After the first few test methods the user will probably be fed up with constantly having to manually introduce test code in so many different text boxes without any handy auto complete feature like in modern IDEs. The tool also lacks some relevant features, like the possibility of creating tests that expect a given exception to be thrown, or to generalise on a generic data type.

To apply REASSERT, changes were introduced to methods of the target of our study, com. puppycrawl.tools.checkstyle.api. After introducing the changes, the tests started failing, as expected. However, using REASSERT it was possible to fix them, using the new functionality that was used, by altering the assertions of the test case.

Finally, JWALK was successfully applied to a class of the case study, the FastStack class. This had to undergo some changes in order for JWALK to properly test it and produce accurate test oracles. The problem was that some code used in the class was not quite deterministic and was causing some problems for JWALK.

```
public class FastStack<E> implements Iterable<E> {
    ...
    /*@ requires this != null; */
    /*@ requires this.mEntries != null; */
    /*@ ensures this.mEntries != null; */
    /*@ ensures \result == 0 || \result == 1 || \result == 100; */
    /*@ ensures \typeof(this.mEntries) == \old(\typeof(this.mEntries)); */
    public int size() {
        return mEntries.size();
    }
    ...
}
```

Listing 3.14: Example of incomplete invariants detected by DAIKON.

3.3.3 Tools Not Tested

The tools that were not applied with success to <code>com.puppycrawl.tools.checkstyle.api</code> were: JM-LUNITNG, JAVA PATHFINDER, UNITCHECK, UDITA, ECLAT and PEX. In this section we discuss why this happened.

The first tool that was not able to run successfully on the subject of this study was JMLUNITNG, the problem arising from the classes not being able compile under the jml4c compiler, a compiler for JML specifications, which is a requirement for JMLUNITNG. The tool was applied to an example annotated class provided at jml4c's website. This new class consists of a simple representation of a bank account and JML is used to specify legal operations over it. Executing the generated test class without providing any additional inputs the tool was able to generate 75 test cases of which 51 were skipped and 5 failed. It is worth mentioning that the tests that failed were due to the incorrect handling of null values in the constructor and main methods of the class. The test cases skipped are due to the inputs not meeting the specified pre-condition. Some more values were provided to the test cases through strategy files generated by the tool, which was then able to generate a total of 588 test cases of which 402 were skipped and the same 5 failed, as is illustrated in Table 3.22.

	Total	Passed	Skipped	Failed	Used Tests (%)	Unused Tests (%)
Without custom values	75	19	51	5	25.33	74.67
With custom val- ues	588	181	402	5	30.78	69.22

Table 3.22: Comparison of JMLUnitNG results.

Continuing to JAVA PATHFINDER and related tools, UNITCHECK and UDITA, it was not possible to successfully apply them to com.puppycrawl.tools.checkstyle.api, nor Checktyle, due to the dependency of Checkstyle on Java Native interface methods. A discussion was opened with the team behind JAVA PATHFINDER, to learn a way to work around this limitation. It was concluded that making Checkstyle work with JAVA PATHFINDER would consume too much time and there would be little or no return from it. In addition, UNITCHECK is very unstable at the time and UDITA is not working properly.

ECLAT was unable to run due the lack of maintenance of the project. Dependencies were broken and the tool will not run regardless of the target system. PEX was not applied, because it is meant for systems implemented with the .NET framework.

KORAT was not successfully applied to Checkstyle. The results achieved with KORAT were either too loose or too strict, and proper ones were never generated. However, it was applied to some examples provided by their respective developers.

3.4 Summary

The tools covered in this case study are regarded as the ones that could aid in the testing phase of the software development cycle. Other tools exist which either were not available for public use or there is no usable implementation available. These include TestEra [Marinov and Khurshid, 2001], Rostra [Xie et al., 2004], Symstra [Xie et al., 2005] and JTestCraft [den Hollander, 2010].

Some commercial tools were also found, like JTest [Parasoft, 2011] and Agitar [Agitar Technologies, 2011], though very little is known of how these solutions work. However, according to Boshernitsan et al. [2006], Agitar makes use of the the DAIKON invariant detector to generate invariants and then uses them to improve the test generation process. At the time of writing there is no pricing information for these tools in their homepage. Evaluations are available for enterprises.

Other tools that were not considered for this case study are tools that rely on the existence of models for the tool under test, like TGV [Jard and Jéron, 2005] and TorX [van Osch, 2005]. The reason these tools were discarded was due to the common lack of a model from which to extract test cases, like Checkstyle, or in some cases the existence of an outdated model.

Chapter 4

RAIKON - An Experiment in TCG Tool Interoperability

There are two ways to contribute to the advance of technology - either by inventing something new and useful, or by showing that the current state-of-the-art tools can be put together in a way such that the whole is better than the sum of its parts.

Tool integration, the second alternative above, is a challenging activity. We follow its path in building a new TCG tool - RAIKON - by putting together two tools reviewed in the previous chapters - RANDOOP and DAIKON. The inspiration came from Eclat, a tool that attempts to infer an operational model of the target software from existing test cases and then applies random testing to check the inferred model . Our tool is named RAIKON, a construction from RANDOOP and DAIKON. RAIKON aims to be a replacement to the abandoned Eclat, while introducing some new functionality.

Like ECLAT, RAIKON attempts to infer a model of correct behaviour of a system from an existing test suite using DAIKON. Then, it uses RANDOOP to generate assertions that check the model, achieving some degree of acceptance testing. The new assertions go beyond the ones that are generated by default by RANDOOP, which are meant for regression testing. The most interesting feature introduced over ECLAT is the possibility to generate input values based on the inferred model. Figure 4.1 depicts the overall architecture of RAIKON.



Figure 4.1: Architecture for RAIKON

RAIKON should improve upon the input domain coverage and help developers in creating complete test suites. RAIKON should also be able to find more flaws in faulty software, since failing tests may be violating the inferred model and that may mean that there is an error in the software

CHAPTER 4. RAIKON - AN EXPERIMENT IN TCG TOOL INTEROPERABILITY

or that the inferred model is incomplete. This is one of the reasons not to completely trust RAIKON, as the model may be incomplete. Ideally RAIKON would use a user provided model of operation. Inferring the model from an existing test suite is meant as a backup solution.

4.1 Implementation

To implement the core of RAIKON, first we have studied how RANDOOP implements checks and how the checks are added to sequences. We also had to study how DAIKON handles invariants. Starting with DAIKON, we have learned that the invariants are divided as unary, binary and ternary according to the number of variables involved in it. After analysing RANDOOP, we have learned that there are no classes implementing generic *n*-ary checks. It became clear that it would be necessary to implement a similar architecture for the checks in RAIKON. So it was decided to implement one class for each arity: unary, binary and ternary. To simplify the conversion between invariants from DAIKON to checks of RANDOOP, a field was added to the newly implemented classes to indicate what operation it intends to express. This allowed a reduction of the number of classes needed to convert the invariants into checks. Finally, to glue the whole process, a sequence visitor (which is how RANDOOP finds which checks should be added to a given statement) was implemented to convert invariants from DAIKON into checks of RANDOOP on a statement-by-statement approach.

In order to support pre-state invariants inferred by DAIKON, the sequence visitor was extended to create a deep copy of the object on which the statement is being invoked. The deep copy is done with a small library named JAVA DEEP-CLONING LIBRARY¹. This library relies on standard reflection mechanisms. This decision introduced a limitation in the tool. When the object to be cloned is too big and complex, i.e. involving too many nested objects, the sequence generation process is delayed.

Since DAIKON refers to parameters by name, and since parameter names are not available by reflection, it was necessary to use a byte code engineering library to gather that information. DAIKON relies on the BCEL² library. However, using BCEL proved not easy to use, and we eventually opted for another such library, PARANAMER³, which provides a much simpler way to get parameter names. This decision has limited the tool to only being able to get parameter names for classes whose debugging information is available, meaning that the classes need to be compiled with the debug flag on.

To generate input values from the inferred entry conditions, it would be necessary to use a constraint solving library to find all the values in range of the bounds defined by the conditions. Two major libraries were found in this respect, $JACOP^4$ and $CHOCO^5$. Experiments were made with both solvers and JACOP was the chosen as the solver to include in RAIKON. In our experiments with both solvers, we have learned that CHOCO was not as reliable and not as scalable as JACOP and that both solvers however lack the support for floating point numbers.

Some pragmatic choices were also made in order to make for some missing functionality in DAIKON. For instance, there is no simple way of getting the name of a variable without tokens introduced by DAIKON. Therefore, it was decided to apply regular expression recognition to the variable names to eliminate each token. This introduces some delay while processing the correct

¹Available at http://code.google.com/p/cloning/

²Available at http://commons.apache.org/bcel/

³Available at http://paranamer.codehaus.org/

⁴Available at http://www.jacop.eu/

⁵Available at http://www.emn.fr/z-info/choco-solver/

variable name. The internal API of DAIKON also has some limitations when attempting to learn what the type of a given variable is. Some methods for such a purpose are incomplete, others non existing, so it was decided to attempt to complete such methods as much as possible.

Currently, RAIKON is available at http://code.google.com/p/raikon/. At the moment it is only available as a command-line tool. All the usual RANDOOP options are still available, new options being added to control the extra functionality. The options are the following:

use-daikon: boolean Signals that this run will use DAIKON artefacts (Slower) [defaults to false].

invoke-daikon:boolean Signals that this run will invoke DAIKON to produce an invariants file (Slower) [defaults to false].

test-case: string The fully-qualified name of a test case.

- projectPrefix:string The project prefix to filter out unwanted classes. Includes only classes
 with the prefix.
- exclusionPattern Complements projectPrefix. Excludes specified classes.
- **heapSize:string** The heap size for invoking DAIKON. Formatted as for the JVM [defaults to "512m"].

invFileName: string The name of the file that will store the invariants [defaults to "invariants"].

disable-regression:boolean Signals that this run will disable regression checks [defaults to false].

daikon-config:string Specifies the DAIKON configuration file, if any.

lowBound:int Lower bound for the input search domain [defaults to -100].

uppBound:int Upper bound for the input search domain [defaults to 100].

maxSolutions: int Maximum number of solutions [defaults to 100].

4.2 Evaluation Set Up

To evaluate RAIKON, three experiments were conducted. The first one addressed FASTSTACK class of the subject of our tool review, Checkstyle. The other two experiments were conducted on the Apache Commons Collection⁶, one on an arbitrary package of the project and the other on the whole project. The reason for choosing Commons Collections is due to RANDOOP having been evaluated on it with success [Pacheco, 2009]. These three experiments also allowed us to evaluate RAIKON on three different levels, a single class, a package and a whole system.

We will measure the number of generated test cases and the input bound coverage, which is measured by the tool itself. We measure line and method coverage to compare both tools effectiveness in covering code elements. They are measured with the EclEmma⁷ plugin for Eclipse. We also measure a ratio of assertions by cyclomatic complexity, to achieve a better idea of the amount

⁶Available at http://commons.apache.org/collections/

⁷Available at http://www.eclemma.org/

CHAPTER 4. RAIKON - AN EXPERIMENT IN TCG TOOL INTEROPERABILITY

of assertions generated. Assertions are measured with an <code>egrep</code> script and cyclomatic complexity is measured with JavaNCSS⁸. Finally, we also compare the number of generated failing test cases, measured with JUnit.

To perform the experiments we ran both tools, RANDOOP and RAIKON, in increasing time spans, starting at 0 and incrementing 5 seconds after each run. We repeat this process until the number of generated test cases starts to stabilise, or the time reaches 900 seconds. In between runs we take note of the desired metrics. The sequences generated have a maximum size of 5 statements. The flow chart in Figure 4.2 illustrates the steps for the experiments, the results of which are detailed in the following sections.



Figure 4.2: Flow of the experiments.

4.3 Results

4.3.1 Experiment A: Single Class

This experiment started with the most simple case where RAIKON was used, the FastStack class from com.puppycrawl.tools.checkstyle.api. Both tools were ran up to 120 seconds. Figure 4.3 provides a comparison of results over the elapsed time. RAIKON was able to generate more sequences, i.e. more test cases in a 120 seconds time span, RANDOOP was able to generate approximately 6,500 while RAIKON was able to generate around 21,000. RAIKON was also able to significantly increase the input coverage in comparison to RANDOOP. RAIKON managed to cover around 7,700 distinct inputs while RANDOOP only managed to cover about 200. Both tools managed to cover all the lines and methods of the FastStack class. When comparing the ratio of number of assertions by cyclomatic complexity, RAIKON also came ahead of RANDOOP achieving a ratio of almost 15,000 against the ratio of 1,000 achieved by RANDOOP. When comparing the number of failed tests RANDOOP did not generate any failing test cases while RAIKON generated around 15,000 failing test cases. Running RAIKON adds an overhead of about one second to read the invariant file and to generate the new inputs.

⁸Available at http://www.kclee.de/clemens/java/javancss/



Figure 4.3: Results for FastStack

4.3.2 Experiment B: Package

The package from Commons Collections chosen to perform our second experiment was, org. apache.commons.collections.collection which holds classes that extend the standard Java interface java.util.Collection. Results for the experiment can be found in Figure 4.4. This experiment took about 240 seconds to stabilise. The results for this package are also very encouraging. RAIKON generated around 22,750 test cases while RANDOOP generated 22,000. Still, the number of covered inputs was significantly increased with RAIKON managing to cover 8,672 distinct inputs and RANDOOP covering 2,781. Both tools achieved similar results in line and method coverage. The line coverage came close to 42% and method coverage was about 40%. When comparing the ratio of assertions by cyclomatic complexity, RAIKON was also able to surpass RANDOOP with a ratio of around 1,400 against the 400 achieved by RANDOOP. RAIKON generated close to 15,500 failing test cases while RANDOOP generated around 1,000 failing test cases. RAIKON took around two seconds in reading the invariants file and generating new inputs.

4.3.3 Experiment C: Whole System

Our experiment with the whole Commons Collections library, was not as successful. This experiment was conducted to a limit of 900 seconds after which only RANDOOP was able to stabilise. RANDOOP was able to generate around 360,000 test cases achieving a coverage of 1,450,000 inputs while RAIKON was only able to generate about 123,000 test cases, that cover around 410,000 inputs. In terms of coverage both tools were close to each other. However, RANDOOP came out on top with a line coverage of 41.9% and method coverage of 60%, while RAIKON managed a line coverage of 39.7% and method coverage of 57.4%. Comparing the rate of assertions by cyclomatic complexity, RANDOOP and RAIKON were equal up until 600 seconds, where RANDOOP stopped generating assertions. RAIKON achieved a rate of 147.41 assertions by cyclomatic complexity and RANDOOP achieved a rate of 105.50. RAIKON also generated more failing tests than RANDOOP with a total of 21,040 failing tests and RANDOOP generating 3,940 failing tests. Figure 4.5 Illustrates the results obtained over time.

4.4 Discussion

We can see that the metrics measured for both tools increase proportionally over time. It can also be seen that the generation of test cases starts growing at a lower pace after a given moment. This is due to the tools generating unique test cases, so the number of generated tests decreases over time. When the tools test all methods, they permute the sequences in order to generate new ones.

Overall, RAIKON preformed better than RANDOOP. In the smaller experiments RAIKON was able to surpass RANDOOP. It significantly exceeded the coverage over input domain, while generating more test cases. The largest experiment, with the whole library, was less successful. RANDOOP was able to generate more test cases and to cover more inputs. This is due to RAIKON's overhead when visiting sequences. Nonetheless, when comparing the number of failing tests, RAIKON surpassed RANDOOP. This could mean that RAIKON has more potential to find more flaws in the software under test.

Over time, we believe that RAIKON can preform better than RANDOOP on all scenarios, when its limitations are lessened.


 $Figure \ 4.4: Results \ for \ {\tt org.apache.commons.collections.collection}$





(c) Comparison of covered lines.



(e) Comparison of assertions by cyclomatic complexity (142).

1,500,000 1,125,000 750,000 375,000 0 0 0 150 300 450 600 750 900 Time (in seconds) • RANDOOP • RANDOOP • RANDOOP

(b) Comparison of covered inputs.



(d) Comparison of covered methods.



(i) companson or raining too

Figure 4.5: Results for Commons Collections

Chapter 5

Conclusions

This is the final chapter of this dissertation which sums up the whole work, presents answers to the research questions and finally, it opens up avenues for future research.

Summing up, we have conducted two reviews: a literature review, where literature of the prominent topic, software testing, of our work was reviewed and a tool review, where we have evaluated existing tools. Then we gave our own contribution to the field by showing that, rather than inventing new tools for TCG one benefits from combining existing tools so as to get the best of them. Our choice was to merge RANDOOP with DAIKON, leading to RAIKON, whose implementation and evaluation was performed.

Our literature review allowed us to learn more about TCG and related techniques in a theoretical approach. We have also learned of the major challenges in testing. We have studied random testing, bounded exhaustive testing, mutation-driven testing and model-based testing. These techniques have their advantages and disadvantages, but they all aim at reducing the cost of testing in a development environment. We have also learned the current state of the art in constraint inference/solving techniques, since they are very closely related with software testing. The most relevant uses of constraint inference techniques were found to be: invariant detection, symbolic execution, dynamic analysis and concolic execution. While the constraint solving techniques studied amounted to TCG satisfiability solving. This study allows to answer our first research question, *What methods are there for Java unit test case generation?*. Our research has revealed that the most common methods of for the Java language are random testing, bounded exhaustive testing and model-based testing. We have also learned about mutation-driven testing, but this is a recent technique and no public implementation of this technique was available.

Finally, our literature review has shown some techniques that could either directly or indirectly useful to TCG, namely invariant detection and input classification. Invariant detection is a technique that relies on constraint inference to infer invariants from source code and input classification is a technique that classifies inputs based on their effects on the target software.

Our tool review, compared many of the tools uncovered in our literature review. Not all of the tools were successfully applied to our study, Checkstyle. Nonetheless, it was shown that the tools that were not applied, were either inappropriate for the tool under study or simply unstable at the time of the case study. The most interesting tools were RANDOOP, FEED4JUNIT and JMLUNITNG. RANDOOP is a very interesting tool that automatically generates a regression testing suite. This tool captures program executions and builds test cases from the observable states of the software. It is however, incapable of exploring a significant part of the software mostly due to its random input generation strategy, which lacks the capability to generate more complex inputs. FEED4JUNIT is very similar to both QUICKCHECK for Java and JCHECK. It does however, provide some additional features that make it more appealing that the other two. Like the possibility to feed test data through a csv file.

Even though JMLUNITNG was not successfully applied to the subject of our study, it was found to

CHAPTER 5. CONCLUSIONS

be very interesting by being able to extract oracles from JML specifications. It also features some reflection mechanisms helping to improve test data generation. Nonetheless, it allows users to provide their own input data. The major set back of JMLUNITNG is that it requires JML specifications. Another problem of JMLUNITNG is that it generates too many unused test cases, leading to a considerable overhead in TCG.

In our tool review we have also found some very interesting tools. The ones that we would like to point out are DAIKON and REASSERT. DAIKON is able to take some of the already written test cases and infer code invariants. This is an interesting functionality, which we found could be used by other tools to improve TCG, as was shown with RAIKON. REASSERT is another interesting tool to help developers in fixing broken test cases.

Let us raise some considerations about the tool review. Firstly, it was noted that some of the tools covered by our study were not suited for the case study, Checkstyle. However, it was decided to continue using Checkstyle as case study due to the previously stated reasons, i.e. for being a common Java application with a history in software engineering research [Cornelissen and Moonen, 2007; Zaidman et al., 2007, 2008; Lubsen et al., 2009]. Another relevant observation about the tools that were covered by this study is that they are far more usable when they come integrated in an IDE, or some form to apply the tool through a GUI. Command line tools are, most of the times, quite complicated to set up, especially in real life applications whose class paths are not always easy to configure.

A last consideration about the studied tools, is that they do not seem ready to be used by the software industry. Some tools, like QUICKCHECK for Java and FEED4JUNIT, even though capable of significantly improving test cases by increasing their input space coverage, are still not automatic enough to efficiently reduce the test case creation cost. In addition, development and maintenance of these tools appear to be fairly slow. Other tools like JAVA PATHFINDER and UDITA are more of research tools than proper tools for use in industrial environments, as their set up is not always simple and in some cases there is very little or no introductory documentation available. However, due to their nature these tools are in constant development. Finally, this study covers some tools that are both stable and easy to use, as is the case of RANDOOP and DAIKON, but are also usable in research contexts by allowing new approaches to solve the problems they aim at solving.

Tables 5.1 and 5.2 sum up all information about the tools covered by the review, enabling us to answer the second research question which triggered the work. *How do the existing methods compare to each other?* The answer is that there are several tools available for TCG. However, some tools are widely unstable or unable to efficiently reduce the costs of software testing.

Concerning the last research question, *Is there an opportunity to integrate such a method in a tool for fully automated test suite augmentation?*, we have implemented, RAIKON, a tool that integrates RANDOOP and DAIKON. The tool uses an operational model inferred by DAIKON to improve the assertions used by RANDOOP and to generate a higher number of valid inputs. This effectively increases the coverage over the input domain and in some cases improves the line coverage in comparison with the test suites generated by RANDOOP. However, in larger systems, RAIKON works significantly slower in comparison with RANDOOP, but that could be due to some limitations of RAIKON, which can be improved. Nonetheless, the results obtained from RAIKON were good enough to inspire future work on the tool. Which answers our last research question affirmatively.

There are, however, some future improvements to RAIKON. First and foremost, the tool limitations should be amended. In an attempt reduce the delay introduced by the deep copy strategy there is a possibility of using serialisation mechanisms instead of reflection. To get the parameter names of a statement, a new strategy also needs to be implemented. For example using the same library as DAIKON, BCEL. Or by extending DAIKON to also provide the parameter position in

	Randoop	QUICK	CHECK	JCHEC	¥	Feed4	JUNIT	T2		JMLUNITNG	UDITA	Рех
Language	Java & .NET	Java		Java		Java		Java		Java	Java	ŃЕТ
Price	Free	Free		Free		Free		Free		Free	Free	Free
Level of Tests	Unit	Unit		Unit		Unit		Unit		Unit	Unit	Unit
Scope	TCG	PUT work	frame-	PUT work	frame-	PUT work	frame-	PUT fi work	'ame-	TCG	Input / Oracle Generator	TCG
Adequacy Criterion								/ Spec.	Valid.	Spec. Valid.		/ Spec. Valid.
Incremental	No	No		No		No		No		No	No	No
Previously Generated Tests	Overwrites	Not able	applic-	Not able	applic-	Not able	applic-	Regress tests	ion	Not applic- able	1	Overwrites
Integration With Manual Tests	Co-Exist	Co-Exi	st	Co-Exi	st	Co-Exi	ist	Co-Exist		Co-Exist	Co-Exist	Co-Exist
Integration With De- velopment Tools	Eclipse	None		None		None		None		None	None	Visual Studio
Code or Bin- ary Level	Binary	Code		Code		Code		Code		Source	Binary	Binary
Configuration Effort	Low	Low		Low		Low		Low		Median	Impossible to say	Low

	Java Pathfinder	UNITCHECK	Korat	Reassert	Daikon	Auto JUnit
Language	Java	Java	Java	Java	Java	Java
Price	Free	Free	Free	Free	Free	Free
Level of Tests	System	Unit	Unit	Unit	Unit	Unit
Scope	Bytecode Model Checker	Unit Model Checker	Input Generator	Test Maintenance	Invariant Detection	Test Stub Gener- ator
Adequacy Cri- terion	Prop. Valid. /	Prop. Valid. /				Branch Coverage
Incremental	No	No	No		Yes	Mo
Previously Gener- ated Tests						No Integration
Integration With Manual Tests	Co-Exist	Co-Exist	Mo	Possible	Can Help	No
Integration With Development Tools	Eclipse / Netbeans	Eclipse / Ant	None	Eclipse	None	None
Code or Binary Level	Binary	Binary	Binary	Code	Binary	Code
Configuration Effort	Varies with the	Low	Low	Low	Low	Low
		Table 5.2	2: Interesting Tools :	Summary.		

CHAPTER 5. CONCLUSIONS

the parameter array, thus dispensing the need for both PARANAMER and BCEL. To patch up the input generation limitation, of only being able to generate integer values, the solver, JaCoP could be extended to support floating point numbers, or we could implement a similar library, whose focus would be the generation of input data based on the conditions inferred by DAIKON. Also, the improvement of the internal API of DAIKON should somewhat improve performance and simplify some internal logic of RAIKON.

Other than repairing the discussed limitations, RAIKON could greatly benefit of more integration with other development tools. Like the Eclipse IDE or the Maven project manager. Other improvements that could benefit RAIKON are improvements on both of the tools that make RAIKON. A suggestion to improve DAIKON would be use existing specifications when inferring invariants. Instead of relying solely in execution traces, DAIKON could also attempt to use existing specifications, such as JML specifications or existing abstract models of the software, to enrich the one it infers. RANDOOP could offer more control on the sequences it generates, allowing a differentiation of *set up* statements, i.e. statements that prepare the instance, from statements that test the instance. Also, in the future, RANDOOP could benefit from a combinatorial testing strategy, in which all possible input combinations are tested. Table 5.3 sums up a profile of RAIKON.

Language	Java
Price	Free under the MIT licence (see MIT [2011])
Level of Tests	Unit level
Scope	Test case generator
Adequacy Criterion	Fault detection / specification coverage
Incremental	Existing test cases are used to infer a model of correct operation
Previously Generated Tests	Overwrites or complements previously generated test cases
Integration With Manual Tests	Co-exist with manual tests
Integration With Development Tools	Only available as a command line tools
Code or Binary Level	Works at the binary level
Configuration Effort	The standalone command line tool requires a proper setup of the class path variable
Quality of Produced Tests	Produces complete test cases, with inputs and test or- acles. Some of the produced tests seem to contain prob- lems that prevent the test classes from compiling and therefore running. Nonetheless, they are easy to fix and execute

Table 5.3: RAIKON Summary.

In conclusion, our work has opened some paths for future research. There is already a wide variety of TCG tools whose source code is open and in some cases, there is room for improvement. For instance, integrating KORAT with FEED4JUNIT in order to allow the generation of more complex

CHAPTER 5. CONCLUSIONS

data types to pass them on to a PUT. DAIKON could be combined with KORAT, in order to produce a valid predicate for the structure to be generated from the invariants inferred by DAIKON. Another avenue for future research is the generation of more abstract test oracles. The invariants inferred by DAIKON are influenced by execution traces, and when the number of execution traces is not enough, the invariants inferred will reflect the execution traces, thus the invariants will not reveal appropriate properties for acceptance testing. Allowing DAIKON to gather information from other means than executions might improve the inferred invariants. These experiments could produce advancements in TCG.

Summing up, tool interoperability shows potential [Sun et al., 2009; Marín et al., 2011]. The effort put in combining tools is payed off by the results, which in most cases, are better than the sum of its parts. This dissertation is proof of that, as our work showed with RAIKON.

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Index of Terms

acceptance testing

a test conducted to determine if the requirements of a specification or contract are met. 6, 7, 51, 64

adequacy criterion

measures the quality of the test suite following a given criterion. 7-9, 19, 73

assertion density

adequacy criterion that measures number of assertions per line of code. 7, 8

black-box testing

tests to validate a specification without any knowledge of the code. xi, 5, 6, 74

bounded exhaustive testing

testing technique that exhaustively tests the input domain. 9, 10, 12, 14, 18, 42, 44, 59

concolic execution

a technique that results from the combination of symbolic execution and dynamic analysis. 15, 16, 18, 35, 59

constraint inference

techniques to infer simpler constraint systems from more complex ones. 8, 9, 15, 16, 59

constraint satisfaction problems

mathematical problems defined as a set of objects whose state must satisfy a number of constraints. 15

constraint solving

techniques to solve constraint systems. 9, 15, 16, 52, 59

cyclomatic complexity

metric that measures the number of linearly independent control flow paths. 53-58, 78-97

dynamic analysis

techniques that attempt to infer constraint systems from actual executions. 9, 15, 16, 59, 73

fault detection

adequacy criteria that measure defect detection effectiveness. 7, 61, 74

formal concept analysis

a way of automatically deriving an ontology from a map of objects and their properties. 9, 10, 99

formal software verification

a process of formalising software validation. xi, 3, 76

functional tests

type of tests that verify that a is abiding a design specification. 7

genetic algorithms

group of algorithms that search for solutions by evolving existent candidates. 12

grey-box testing

a method that combines black-box testing and white-box testing, validates a certain specification using knowledge of internal algorithms and structures. xi, 5, 6

input bound coverage

measure of how much of the input is covered. 7, 12, 23, 37, 46, 51, 53, 60-62

input classification

technique to classify inputs by order of criticality. 17, 59

integration testing

tests for modules or sets of units. 6

invariant detection

research field that aims to detect invariants from source code. 16, 17, 59

model checking

technique to formally validate software. 3, 31

model-based testing

technique that generates test cases from models. 9, 10, 12, 14, 18, 59

mutation

artificial change introduced in the source code, with the aim of detecting it. 7, 12, 14

mutation analysis

technique to measure how effective a test suite is in fault detection. 7, 12

mutation-driven testing

Test case generation technique to detect mutants. 9, 10, 12, 14, 59

non-functional tests

tests for non-functional requirements, such as performance, stability, etc. 6

ontology

a formal representation of knowledge as a set of concepts within a domain. 9

parameterised unit testing

generalisation of unit testing allowing inputs as parameters. xi, 6, 12, 23, 25, 27, 37, 46, 61, 64

random testing

Test case generation technique that randomly generates test cases. 9, 10, 12, 14, 18, 27, 51, 59

regression testing

a type of testing that seeks to uncover new errors in existing functionality after making changes. 6, 7, 23, 51, 59

satisfiability solver

determine if the variables of a given boolean formula can be assigned. 16, 59

software

set of functions that compose an application. 3–6, 8, 9, 12, 14–19, 23, 31, 32, 35, 40, 45, 50, 51, 59, 60, 74, 75

software bug

flaw that causes incorrect behaviour of code. 10, 12

software testing

process intended to check the quality and correctness of software. 3, 5–10, 12, 14, 15, 17–19, 21, 27, 37, 40, 42, 44, 50, 59, 60

specification mining

techniques that attempt to generate program specifications from executions. 17

stakeholder

a person with an interest or concern in a product. 6, 7, 18

structural coverage

coverage criteria that measure how many code elements are executed. 7, 29, 46, 53, 54, 56, 60–62

symbolic execution

techniques that attempt to generate a constraint system through static execution. 9, 15, 16, 59, 73

system testing

type of tests that check if a given software is behaving as expected. 6, 7, 35

system under test

application or component being targeted for testing. 6, 12, 14, 16, 17, 35, 62

test

sequence of instructions that assert a given expected behaviour. 4–8, 10, 12, 14, 17–19, 23, 27, 35, 42, 46, 48, 51, 54, 59, 60, 76

test case generation

techniques that automatically generate test cases. vii, xi, 4, 8–10, 12, 14–18, 21, 23, 27, 29, 31, 37, 40, 45, 46, 51, 59–61, 63, 64, 74–76

test data generator

component of a test case generation that generates data. 9, 17, 23, 60

test oracle

component of a test case generation that asserts expected behaviour. 6, 12, 14, 15, 18, 23, 35, 60, 64

test suite

a set of test cases. 3, 4, 6-9, 12, 17, 19, 51, 52, 60, 73

theorem proving

a formal software verification technique that proves program correctness. 3

unit testing

testing of small application units. 6, 35, 74

white-box testing

tests to check internal workings and structures. xi, 5, 6, 74

Appendix A

Detailed Results

The tables that lead to	Figure 4.3.	Figure 4.4.	Figure 4.5 c	an be found bellow.

Time Reading Invariants	0:00:00.738
Time Solving Constraints	0:00:00.260
Total Time	0:00:00.998

Table A.1: Set up time for FastStack

Time Reading Invariants	0:00:01.502
Time Solving Constraints	0:00:00.647
Total Time	0:00:02.149

Table A.2: Set up time for <code>org.apache.commons.collections.collection</code>

Time Reading Invariants	0:02:01.674
Time Solving Constraints	0:00:03.270
Total Time	0:02:04.944

Table A.3: Set up time for Commons Collections

	Gene Sequ	erated ences	_ C	vered nputs	Line Co	verage (%)	Method C	overage (%)	L Cyc Comple	Asserts/ lomatic xity (14)		uiling 「ests
Time	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	RAIKON
0	0	0	0	0	0.0	0.0	0.0	0.0	0	0	0	0
თ	2,266	1,783	184	2,240	100.0	100.0	100.0	100.0	376.14	1,028.21	0	886
10	3,001	3,817	187	4,085	100.0	100.0	100.0	100.0	483.64	2,405.93	0	2,337
15	3,524	5,391	187	4,269	100.0	100.0	100.0	100.0	552.79	3,701.79	0	3,660
20	3,875	7,163	187	4,606	100.0	100.0	100.0	100.0	601.64	4,784.21	0	4,777
25	4,082	8,438	189	5,207	100.0	100.0	100.0	100.0	635.79	5,722.50	0	5,751
30	4,425	9,419	189	5,315	100.0	100.0	100.0	100.0	684.79	6,657.07	0	6,728
35	4,655	10,437	189	5,955	100.0	100.0	100.0	100.0	698.36	7,260.07	0	7,558
40	4,917	11,095	189	6,165	100.0	100.0	100.0	100.0	744.00	8,119.21	0	8,255
45	5,095	11,880	189	6,424	100.0	100.0	100.0	100.0	767.29	8,633.43	0	8,804
50	5,241	12,664	189	6,643	100.0	100.0	100.0	100.0	804.21	9,285.93	0	9,495
55	5,443	13,156	189	6,794	100.0	100.0	100.0	100.0	823.71	9,765.93	0	9,998

	Gen Sequ	erated Jences	= O C	vered nputs	Line Cc	vverage (%)	Method C	toverage (%)	Cy Comple	Asserts/ /clomatic exity (14)	Ë.	ailing Fests
Time	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon
60	5,599	14,244	189	6,970	100.0	100.0	100.0	100.0	854.21	10,158.50	0	10,420
65	5,627	14,779	191	7,051	100.0	100.0	100.0	100.0	864.93	10,604.79	0	10,879
70	5,695	15,429	191	7,182	100.0	100.0	100.0	100.0	874.00	11,308.79	0	11,625
75	5,781	16,131	191	7,280	100.0	100.0	100.0	100.0	890.93	11,885.00	0	12,234
80	5,944	16,336	191	7,316	100.0	100.0	100.0	100.0	904.00	12,063.00	0	12,422
85	6,107	16,925	191	7,387	100.0	100.0	100.0	100.0	919.29	12,686.86	0	13,076
06	6,171	17,242	193	7,415	100.0	100.0	100.0	100.0	934.64	13,018.21	0	13,427
95	6,248	17,897	193	7,487	100.0	100.0	100.0	100.0	945.71	13,521.64	0	13,954
100	6,305	18,488	193	7,552	100.0	100.0	100.0	100.0	952.79	13,832.21	0	14,287
105	6,397	18,684	193	7,596	100.0	100.0	100.0	100.0	965.29	14,047.79	0	14,521
110	6,452	19,286	193	7,649	100.0	100.0	100.0	100.0	974.36	14,553.43	0	15,060
115	6,521	19,673	193	7,669	100.0	100.0	100.0	100.0	986.36	14,659.29	0	15,174
120	6,571	20,155	193	7,709	100.0	100.0	100.0	100.0	982.43	14,811.07	0	15,338
			Ÿ	able A.5: I	Detailed res	:ults for $\mathbb{F}_{\mathbb{C}}$	astStack ((60-120 s∈	sconds)			

second
(60-120
for FastStack
results
Detailed
Table A.5:

erated lences	L Co	rered Iputs	Line Co	verage (%)	Method C	overage (%)	¢ Cyc Complexi	\sserts/ lomatic ty (142)	. .	uiling 「ests
Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon
0	0	0	0.0	0.0	0.0	0.0	0	0	0	0
1,154	766	1,131	31.1	37.5	30.6	36.7	31.63	58.15	64	617
2,636	1,089	2,253	34.0	38.7	34.7	36.7	73.32	150.80	164	1,597
4,142	1,269	3,071	35.6	39.1	35.7	37.8	115.27	241.37	382	2,537
5,628	1,438	3,762	35.6	39.1	35.7	37.8	157.85	337.18	474	3,537
7,232	1,560	4,433	35.8	39.1	35.7	37.8	197.50	429.25	562	4,514
8,688	1,642	4,844	36.6	39.1	36.7	37.8	234.63	532.39	567	5,604
10,125	1,757	5,354	36.6	39.1	36.7	37.8	267.07	624.35	656	6,577
11,446	1,831	5,784	36.6	39.1	36.7	37.8	292.37	714.97	732	7,547
12,857	1,896	6,118	39.3	39.1	37.8	37.8	324.20	799.34	823	8,424
14,143	2,098	6,530	39.3	39.1	37.8	37.8	346.18	881.18	887	9,300
	2,171	6,823	39.3	39.1	37.8	37.8	363.35	958.64	940	10,122
	0,004	2,004 2,171	2,304 Z,171 0,0Z3	0,004 2,1/1 0,020 09.0	J,304 2,171 0,023 39.3 39.1	ס. יס ו. ישט ט. ישט יו				

Table A.6: Detailed results for org.apache.commons.collections.collection (0-55 seconds)

	Gen Sequ	erated ences	C C	vered puts	Line Co	verage (%)	Method C	overage (%)	Cyc	Asserts/ clomatic ity (142)	Ξ.	illing Tests
Time	RANDOOP	Raikon	RANDOOP	Raikon	RANDOOP	Raikon	Randoop	Raikon	RANDOOP	Raikon	RANDOOP	Raikon
60	20,851	16,543	2,206	7,009	39.3	39.1	37.8	37.8	378.17	1,027.96	988	10,862
65	21,340	17,596	2,330	7,249	39.3	39.1	37.8	37.8	386.13	1,050.78	1,011	10,878
70	21,637	18,552	2,427	7,588	40.5	39.1	38.8	37.8	391.11	1,124.59	1,034	11,899
75	21,841	19,486	2,451	7,808	40.5	39.1	38.8	37.8	394.88	1,216.68	1,051	12,902
80	21,927	20,205	2,582	8,033	40.5	39.1	38.8	37.8	396.30	1,254.15	1,053	13,313
85	21,976	20,764	2,584	8,146	40.5	39.8	38.8	37.8	397.14	1,282.98	1,056	13,637
06	22,024	21,316	2,639	8,292	40.5	39.8	38.8	37.8	397.82	1,313.06	1,058	13,971
95	22,040	21,677	2,653	8,295	41.7	39.8	39.8	37.8	398.15	1,332.27	1,058	14,199
100	22,050	22,022	2,730	8,445	41.7	40.3	39.8	38.8	398.26	1,349.65	1,058	14,392
105	22,051	22,308	2,744	8,495	41.7	40.3	39.8	38.8	398.32	1,359.85	1,058	14,505
110	22,056	22,487	2,755	8,544	41.7	40.3	39.8	38.8	398.35	1,365.96	1,058	14,575
115	22,061	22,622	2,770	8,614	41.7	40.3	39.8	38.8	398.42	1,371.32	1,058	14,633
120	22,065	22,756	2,772	8,651	41.7	40.3	39.8	38.8	398.65	1,378.10	1,058	14,707
		Toble A	7. Dotoilod	07 04 ···000				:				

Table A.7: Detailed results for org.apache.commons.collections.collection (60-120 seconds)

	Gen Sequ	erated ences	_ Co	vered nputs	Line Co	verage (%)	Method C	overage (%)	Cy Complex	Asserts/ clomatic ity (142)	_ . .	ailing Fests
Time	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon
125	22,075	22,853	2,781	8,672	41.7%	40.3%	39.8%	38.8%	398.46	1,378.99	1,058	14,716
130	22,084	23,014	2,821	8,696	41.7%	40.3%	39.8%	38.8%	398.89	1,388.27	1,058	14,814
135	22,089	23,102	2,837	8,720	41.7%	40.3%	39.8%	38.8%	398.81	1,393.61	1,058	14,863
140	22,091	23,187	2,844	8,736	41.7%	41.4%	39.8%	39.8%	398.92	1,397.64	1,058	14,911
145	22,096	23,239	2,861	8,765	41.7%	41.4%	39.8%	39.8%	399.11	1,400.68	1,058	14,944
150	22,101	23,262	2,853	8,773	41.7%	41.4%	39.8%	39.8%	399.02	1,402.04	1,058	14,957
155	22,108	23,314	2,884	8,785	41.7%	41.4%	39.8%	39.8%	399.23	1,404.32	1,058	14,981
160	22,108	23,365	2,885	8,791	41.7%	41.4%	39.8%	39.8%	399.23	1,406.62	1,058	15,006
165	22,110	23,461	2,888	8,799	41.7%	41.4%	39.8%	39.8%	399.26	1,412.30	1,058	15,071
170	22,123	23,462	2,894	8,799	41.7%	41.4%	39.8%	39.8%	399.49	1,412.28	1,058	15,071
175	22,125	23,558	2,897	8,806	41.7%	41.4%	39.8%	39.8%	399.53	1,417.49	1,058	15,130
180	22,125	23,618	2,898	8,816	41.7%	41.4%	39.8%	39.8%	399.53	1,420.46	1,058	15,159
185	22,127	23,648	2,898	8,819	41.7%	41.4%	39.8%	39.8%	399.59	1,421.65	1,058	15,170

Table A.8: Detailed results for org.apache.commons.collections.collection (125-185 seconds)

	Gen Sequ	lerated	= 0 0	vered nputs	Line Co	werage (%)	Method C	overage (%)	Cyc Cyc Complex	Asserts/ clomatic ity (142)	Ца	iling ests
Time	RANDOOP	Raikon	Randoop	Raikon	RANDOOP	Raikon	Randoop	Raikon	RANDOOP	Raikon	Randoop	Raikon
190	22,129	23,659	2,899	8,820	41.7%	41.4%	39.8%	39.8%	399.66	1,423.48	1,058	15,174
195	22,134	23,691	2,914	8,823	41.7%	41.4%	39.8%	39.8%	399.56	1,422.06	1,058	15,187
200	22,139	23,722	2,928	8,825	41.7%	41.4%	39.8%	39.8%	399.74	1,424.58	1,058	15,198
205	22,139	23,758	2,928	8,830	41.7%	41.4%	39.8%	39.8%	399.74	1,425.99	1,058	15,212
210	22,144	23,819	2,930	8,838	41.7%	41.4%	39.8%	39.8%	399.94	1,428.90	1,058	15,236
215	22,150	23,828	2,931	8,839	41.7%	41.4%	39.8%	39.8%	399.82	1,428.49	1,058	15,240
220	22,164	23,886	2,934	8,854	41.7%	41.4%	39.8%	39.8%	400.12	1,430.91	1,058	15,259
225	22,171	23,903	2,935	8,854	41.7%	41.4%	39.8%	39.8%	400.21	1,431.56	1,058	15,265
230	22,174	23,938	2,937	8,859	41.7%	41.4%	39.8%	39.8%	400.25	1,433.06	1,058	15,278
235	22,184	23,971	2,945	8,864	41.7%	41.4%	39.8%	39.8%	400.41	1,434.80	1,058	15,296
240	22,199	24,056	2,966	8,875	42.5%	41.4%	40.8%	39.8%	400.63	1,437.62	1,058	15,322
		Table A.	9: Detailed	results for	l org.apach	e.commons	s.collectio	ns.collect	tion (190-2 4	40 seconds	(\$	

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	Gene Sequ	erated ences		vered nputs	Line Co	verage (%)	Method C	overage (%)	Cyc Complexit	Asserts/ clomatic y (6,848)	Fa	niling Fests
Time	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon
0	0	0	0	0	0.0	0.0	0.0	0.0	0	0	0	0
თ	3,137	780	7,488	1,939	24.6	15.4	38.9	24.1	0.81	0.47	27	88
10	6,049	896	14,557	1,939	28.3	17.8	44.4	28.0	1.05	0.81	62	129
15	9,661	1,010	25,663	5,307	30.4	25.6	47.4	40.2	1.86	1.87	106	295
20	13,585	2,112	37,338	6,000	31.7	26.2	49.1	40.9	2.78	2.26	170	332
25	20,848	3,629	49,403	9,293	32.5	28.3	49.1	43.9	3.68	3.52	216	519
30	24,417	4,197	60,541	10,309	32.8	28.7	50.0	44.4	4.53	4.03	260	587
35	27,979	5,277	73,164	13,386	33.1	29.8	50.6	45.9	5.46	5.10	287	733
40	30,767	5,834	86,797	15,100	33.6	30.3	51.1	46.5	6.49	5.49	334	814
45	34,223	7,051	98,025	19,423	34.0	31.2	51.6	47.7	7.33	6.67	360	1,038
50	36,771	7,470	108,565	20,677	34.2	31.4	51.9	47.9	8.12	7.28	393	1,124
55	39,515	9,136	118,557	24,260	34.5	32.0	52.2	48.8	8.84	8.52	419	1,299

Table A.10: Detailed results for	
Commons C	
ollections (
0-55 seconds)	

	Gen Sequ	erated lences	- 0 0 0	vered nputs	Line Co	verage (%)	Method C	overage (%)	Cy. Complexit	Asserts/ clomatic y (6,848)	Ë,	uiling Tests
Time	RANDOOP	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	RANDOOP	Raikon	RANDOOP	Raikon
60	42,884	10,594	124,008	25,722	34.6	32.0	52.7	48.8	9.24	9.02	435	1,384
65	46,279	10,799	143,930	29,159	34.6	32.9	52.7	49.6	9.27	10.01	438	1,569
70	49,103	12,293	155,477	31,437	35.0	33.1	52.7	49.9	11.57	10.82	527	1,597
75	52,503	12,546	165,331	33,555	35.1	33.1	53.4	50.0	11.76	10.86	531	1,690
80	55,883	14,071	170,882	35,936	35.3	33.5	53.5	50.5	12.71	12.30	565	1,926
85	59,242	15,369	195,100	38,760	35.4	33.5	53.7	50.5	13.52	12.61	601	1,959
06	60,662	15,694	207,554	40,642	35.5	33.9	53.9	50.7	14.43	14.02	628	2,181
95	62,595	17,157	214,439	43,717	35.9	34.2	54.0	51.0	15.97	14.82	696	2,342
100	65,700	17,347	234,944	45,582	35.9	34.2	54.2	51.1	17.09	15.13	732	2,473
105	69,029	17,348	246,313	48,518	36.0	34.4	54.4	51.2	17.91	15.64	760	2,618
110	72,143	18,304	261,937	50,643	36.2	34.4	54.5	51.3	19.02	16.93	770	2,711
115	75,420	18,594	269,775	52,375	36.3	34.5	54.6	51.3	19.38	17.53	806	2,814
120	78,718	18,594	283,644	53,692	36.3	34.6	54.7	51.4	20.69	18.43	852	2,895
							:					

Table A.11: Detailed results for Commons Collections (60-120 seconds)

	Gen Sequ	erated ences	= 0	vered nputs	Line Co	verage (%)	Method C	overage (%)	Cyc Complexity	Asserts/ clomatic / (6,848)	ц П П	viling 「ests
Time	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon
125	81,247	19,007	293,331	55,428	36.6	34.6	54.8	51.4	21.58	19.01	894	2,290
130	85,177	20,601	308,713	58,759	36.9	34.9	54.9	51.7	22.73	21.12	924	3,170
135	86,932	21,906	315,484	62,935	36.9	35.0	55.2	52.0	23.23	22.65	941	3,416
140	88,413	22,182	321,377	63,909	37.0	35.1	55.5	52.1	23.66	23.32	954	3,401
145	88,985	23,129	323,595	66,739	37.1	35.1	55.6	52.2	23.83	24.21	958	3,620
150	89,874	23,843	327,278	68,867	37.1	35.2	55.6	52.2	24.07	25.07	686	3,726
155	90,724	23,846	330,453	68,869	37.2	35.2	55.7	52.3	24.30	25.07	1,014	3,811
160	92,380	24,309	337,234	70,280	37.2	35.2	55.7	52.3	24.77	25.72	1,021	3,811
165	92,840	25,429	338,972	73,587	37.2	35.6	55.8	52.5	24.90	26.99	1,021	3,970
170	98,824	26,441	362,421	76,775	37.2	35.7	55.9	52.6	26.62	28.28	1,123	4,141
175	102,296	27,080	376,274	78,900	37.4	35.8	55.9	52.8	27.63	29.15	1,168	4,256
180	106,172	27,660	391,729	80,766	37.5	35.8	56.0	52.8	28.76	30.09	1,211	4,357
185	108,381	28,777	400,411	84,555	37.6	35.9	56.1	52.9	29.45	31.33	1,251	4,545

Table A.12: Detailed results for Commons Collections (125-185 seconds)

	Gen Sequ	erated iences	- °	vered nputs	Line Co	overage (%)	Method C	overage (%)	Cy Complexit	Asserts/ clomatic y (6,848)	Ë.	illing Tests
Time	RANDOOP	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	RANDOOP	Raikon	RANDOOP	Raikon
190	110,885	29,339	409,989	86,287	37.6	35.9	56.2	53.0	30.13	32.06	1,271	4,662
195	113,079	29,742	418,731	87,634	37.6	36.0	56.3	53.1	30.77	32.60	1,295	4,814
200	113,931	29,836	422,074	87,920	37.7	36.0	56.3	53.1	31.03	32.69	1,297	4,722
205	114,779	30,233	425,441	89,173	37.7	36.1	56.3	53.2	31.28	33.08	1,315	4,738
210	115,313	30,413	427,414	89,679	37.7	36.1	56.4	53.2	31.44	33.26	1,323	4,837
215	116,833	30,414	433,702	89,686	37.8	36.1	56.4	53.2	31.89	33.26	1,341	4,839
220	119,782	30,418	445,671	89,688	37.8	36.1	56.5	53.3	32.74	33.26	1,366	5,101
225	123,173	32,005	459,152	95,017	37.9	36.1	56.5	53.3	33.71	35.13	1,406	5,130
230	123,900	32,159	462,095	95,542	37.9	36.1	56.6	53.3	33.92	35.30	1,406	5,130
235	125,495	33,244	468,451	99,150	38.0	36.1	56.6	53.3	34.38	36.47	1,436	5,306
240	128,458	33,622	479,927	100,399	38.0	36.1	56.6	53.3	35.25	36.87	1,412	5,165
245	143,710	37,398	540,743	112,991	38.3	36.1	56.6	53.7	39.71	41.21	1,742	5,761
250	145,799	37,971	549,433	115,001	38.3	36.4	57.0	53.7	40.31	41.98	1,631	5,966
			Table A.	13: Detail	ed results fo	or Commo	ons Collectio	ons (190-2	50 seconds)			

90-250 second
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Collections
Commons
for
results
Detailed
A.13:
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	Gene Sequ	erated ences	_ 6	vered nputs	Line Co	verage (%)	Method C	overage (%)	Cyc Complexit	y (6,848)	ц П В	uiling 「ests
Time	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	RAIKON
255	148,062	38,274	558,658	116,048	38.3	36.5	57.0	53.8	40.98	42.24	1,656	6,314
260	148,383	38,589	559,951	117,068	38.3	36.5	57.0	53.8	41.08	42.65	1,674	6,162
265	151,561	38,796	572,833	117,822	38.4	36.5	57.0	53.8	42.02	42.91	1,676	6,310
270	152,731	39,465	577,701	119,763	38.4	36.5	57.0	53.8	42.36	43.70	1,705	6,197
275	155,135	40,253	587,689	122,411	38.5	36.6	57.0	53.8	43.08	44.76	1,714	6,463
280	156,307	40,884	592,537	124,547	38.5	36.6	57.0	53.9	43.43	45.53	1,758	6,556
285	159,558	42,371	605,620	129,238	38.5	36.7	57.0	54.0	44.38	47.31	1,790	6,800
290	163,281	42,839	620,745	131,405	38.6	36.7	57.1	54.3	45.48	47.54	1,844	6,937
295	164,431	43,518	625,525	133,113	38.6	36.7	57.2	54.0	45.83	48.61	1,855	7,037
300	168,090	43,583	640,522	133,420	38.7	36.7	57.2	54.1	46.91	48.72	1,891	7,230
305	174,219	44,718	665,667	137,185	38.7	36.8	57.3	54.1	48.75	50.07	1,957	7,299
310	174,762	45,198	667,910	138,916	38.8	36.9	57.3	54.3	48.91	50.44	1,961	7,684
315	175,524	46,426	671,033	143,416	38.8	36.9	57.4	54.0	49.14	51.49	1,972	7,024

Table A.14: Detailed results for Commons Collections (255-315 seconds)

	Gen Sequ	erated lences	- Co	vered nputs	Line Co	werage (%)	Method C	overage (%)	Cy Complexit	Asserts/ clomatic y (6,848)	Ë.	uiling Fests
Time	RANDOOP	Raikon	RANDOOP	Raikon	Randoop	Raikon	Randoop	Raikon	RANDOOP	Raikon	RANDOOP	RAIKON
320	176,768	47,313	676,136	146,150	38.8	36.9	57.5	54.5	49.51	53.33	1,987	7,555
325	180,422	47,915	691,263	148,294	38.8	36.9	57.5	54.3	50.60	54.06	2,022	7,789
330	181,877	48,394	697,528	150,165	38.8	36.9	57.5	54.3	51.04	54.48	2,035	7,872
335	184,499	48,628	708,378	150,897	38.8	36.9	57.5	54.3	51.82	54.84	2,064	7,914
340	186,229	49,387	715,693	153,482	38.9	36.9	57.5	54.4	52.33	55.88	2,075	8,065
345	188,366	50,039	724,512	155,527	39.0	37.0	57.6	54.4	52.96	56.42	2,093	8,158
350	188,913	50,042	726,659	155,537	39.0	37.1	57.6	54.4	53.13	56.42	2,121	8,158
355	190,702	51,711	733,984	160,984	39.0	37.2	57.6	54.6	53.67	57.48	2,104	8,427
360	199,420	53,424	769,785	166,619	39.1	37.3	57.6	54.8	56.26	58.23	2,218	8,727
365	205,825	54,412	796,429	169,975	39.1	37.3	57.6	54.8	58.18	60.26	2,291	8,984
370	207,633	54,751	804,112	171,191	39.2	37.3	57.7	54.8	58.72	61.69	2,306	9,059
375	208,457	55,261	807,409	172,950	39.2	37.3	57.7	54.8	58.97	62.31	2,314	9,062
380	214,953	55,279	833,900	173,015	39.3	37.3	57.8	54.8	60.95	63.03	2,400	8,913
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Table A.15: Detailed results for Commons Collections (320-380 seconds)

	Gen Sequ	erated lences	- 8	nputs	Line Co	verage (%)	Method C	overage (%)	Cyc Complexity	v (6,848)		ailing Tests
Time	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	RANDOOP	Raikon	Randoop	Raikon
385	215,027	56,746	834,250	178,423	39.3	37.4	57.8	55.0	61.34	63.05	2,550	9,604
390	217,998	57,139	846,924	179,633	39.3	37.4	57.8	55.0	61.88	64.87	2,553	9,646
395	220,223	58,211	856,156	183,198	39.4	37.4	57.8	55.0	62.55	65.51	2,386	9,336
400	223,878	58,388	871,203	183,821	39.6	37.4	58.0	55.0	63.66	66.75	2,415	9,401
405	227,876	58,416	888,120	183,895	39.6	37.4	58.0	55.0	64.86	67.00	2,242	9,642
410	229,219	58,884	893,815	185,383	39.6	37.5	58.0	55.1	65.26	67.01	2,474	9,732
415	232,807	59,959	908,862	188,979	39.6	37.5	58.0	55.2	66.36	67.46	2,512	9,910
420	233,220	60,415	910,602	190,594	39.6	37.6	58.0	55.3	66.49	68.74	2,522	9,986
425	233,817	61,531	913,157	194,491	39.6	37.7	58.0	55.3	66.67	69.32	2,584	10,167
430	236,174	61,889	922,722	197,557	39.6	37.7	58.1	55.4	67.39	70.05	2,559	10,163
435	237,765	62,959	929,230	199,329	39.6	37.7	58.1	55.4	67.88	70.55	2,680	10,421
440	244,570	63,235	957,030	200,234	39.7	37.8	58.2	55.4	69.92	72.28	2,691	10,476
445	244,803	64,114	957,961	203,493	39.7	37.8	58.2	55.4	69.99	72.61	2,694	10,623

Table A.16: Detailed results for Commons Collections (385-445 seconds)

	Gen Sequ	erated iences	ö –	vered Inputs	Line Co	verage (%)	Method C	overage (%)	Cycomplexit	Asserts/ clomatic y (6,848)	E E	illing Tests
Time	RANDOOP	Raikon	RANDOOP	Raikon	RANDOOP	Raikon	Randoop	Raikon	Randoop	Raikon	RANDOOP	Raikon
450	248,642	64,826	973,418	205,995	39.8	37.8	58.4	55.4	71.13	73.74	2,737	10,753
455	248,643	65,345	973,422	207,726	39.8	37.8	58.4	55.4	71.13	74.59	2,737	10,842
460	252,403	66,492	989,478	211,809	39.8	37.8	58.4	55.5	72.26	75.19	2,771	11,025
465	252,892	67,078	991,568	213,874	39.9	37.9	58.4	55.5	72.41	76.44	2,805	11,208
470	258,700	67,540	1,016,022	215,543	39.9	37.9	58.4	55.6	74.16	77.13	2,839	11,391
475	259,984	68,176	1,021,425	217,751	39.9	37.9	58.5	55.6	74.54	77.65	2,873	11,574
480	262,807	69,793	1,032,845	223,684	39.9	37.9	58.5	55.7	75.39	78.12	2,910	11,761
485	265,175	70,309	1,042,652	225,374	39.9	37.9	58.5	55.7	76.11	78.48	2,926	11,827
490	266,216	70,492	1,047,283	225,935	39.9	38.0	58.5	55.7	76.42	80.21	2,942	11,893
495	267,770	70,558	1,053,870	226,234	39.9	38.0	58.6	55.7	76.90	80.89	2,958	11,959
500	271,553	71,446	1,070,121	229,422	39.9	38.0	58.6	55.8	78.05	81.01	2,974	12,025
505	275,296	72,674	1,086,331	233,774	39.9	38.0	58.6	55.8	79.17	81.15	2,994	12,091
510	277,314	73,090	1,094,944	235,236	39.9	38.0	58.6	55.8	79.80	81.82	3,019	12,161
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Table A.17: Detailed results for Commons Collections (450-510 seconds)

	Gen Sequ	erated lences	- 8	vered nputs	Line Co	verage (%)	Method C	overage (%)	Cyo Complexit	Asserts/ clomatic y (6,848)		ulling 「ests
Time	Randoop	RAIKON	Randoop	Raikon	Randoop	Raikon	Randoop	RAIKON	Randoop	RAIKON	Randoop	RAIKON
515	283,379	73,743	1,120,118	237,440	39.9	38.0	58.6	55.8	81.63	83.20	3,044	12,231
520	285,701	74,436	1,130,015	241,161	40.0	38.0	58.6	55.7	82.33	83.81	3,069	12,301
525	288,053	74,784	1,140,169	241,345	40.0	38.0	58.7	55.7	83.03	84.54	3,094	12,371
530	288,647	76,519	1,142,513	247,325	40.0	38.0	58.7	55.8	83.22	85.94	3,121	12,443
535	290,262	76,861	1,149,230	248,481	40.0	38.0	58.7	55.9	83.71	87.35	3,149	12,572
540	292,560	78,176	1,158,853	252,923	40.1	38.1	58.7	55.9	84.38	88.01	3,177	12,701
545	295,767	78,193	1,172,399	252,981	40.1	38.2	58.7	55.9	85.36	88.62	3,205	12,830
550	300,035	78,760	1,190,364	254,908	40.2	38.3	58.7	56.0	86.63	90.34	3,233	12,959
555	302,458	79,305	1,200,758	256,763	40.2	38.4	58.7	56.0	87.37	90.35	3,273	13,100
560	305,521	79,872	1,213,669	258,702	40.2	38.5	58.7	56.0	88.30	90.90	3,282	13,211
565	306,325	81,533	1,217,109	264,919	40.2	38.5	58.7	56.1	88.55	91.44	3,291	13,322
570	308,540	81,550	1,226,636	264,987	40.2	38.6	58.7	56.1	89.22	92.09	3,300	13,433
575	309,894	81,555	1,232,400	265,003	40.2	38.6	58.7	56.2	89.63	94.29	3,309	13,544

APPENDIX A. DETAILED RESULTS

Table A.18: Detailed results for Commons Collections (515-575 seconds)
	Gen Sequ	erated lences	<u> </u>	vered	Line Co	verage (%)	Method C	toverage (%)	Cy. Complexit	Asserts/ clomatic y (6,848)	Fa	uiling Fests
Time	RANDOOP	Raikon	RANDOOP	Raikon	Randoop	Raikon	Randoop	Raikon	RANDOOP	Raikon	RANDOOP	Raikon
580	310,391	81,564	1,234,629	265,053	40.2	38.7	58.7	56.2	89.78	94.32	3,311	13,656
585	311,311	81,918	1,238,343	266,273	40.2	38.7	58.8	56.2	90.06	94.32	3,331	13,715
590	314,625	81,998	1,252,440	266,550	40.2	38.7	58.8	56.2	91.08	94.34	3,351	13,774
595	314,864	82,065	1,253,521	266,794	40.2	38.7	58.8	56.2	91.15	94.76	3,371	13,833
600	316,101	82,071	1,258,750	266,801	40.2	38.7	58.8	56.2	91.53	94.92	3,391	13,892
605	317,646	83,204	1,265,543	271,026	40.2	38.7	58.8	56.2	92.00	94.94	3,413	13,952
610	319,895	83,206	1,274,789	271,035	40.4	38.7	59.0	56.2	92.69	94.94	3,430	14,049
615	323,595	84,417	1,290,234	275,547	40.5	38.7	59.1	56.2	93.80	96.35	3,447	14,146
620	326,020	84,932	1,300,696	277,321	40.7	38.8	59.3	56.2	94.54	96.35	3,464	14,243
625	328,398	84,934	1,311,038	278,676	40.9	38.8	59.5	56.2	95.27	97.73	3,481	14,340
630	330,509	86,240	1,319,560	281,849	41.0	38.8	59.6	56.2	95.91	98.21	3,498	14,437
635	330,609	86,584	1,319,978	283,102	41.0	38.8	59.6	56.2	95.94	98.65	3,515	14,562
640	332,355	88,152	1,327,418	288,612	41.0	38.8	59.6	56.3	96.46	99.93	3,532	14,687
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Table A.19: Detailed results for Commons Collections (580-640 seconds)

	Gen Sequ	erated lences	- 8	nputs	Line Co	verage (%)	Method C	overage (%)	Cy Complexit	Asserts/ clomatic y (6,848)	_ 5	uiling 「ests
Time	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon
645	332,860	88,175	1,329,356	288,691	41.0	38.9	59.6	56.3	96.62	100.42	3,549	14,812
650	333,315	89,616	1,331,269	293,647	41.0	38.9	59.6	56.4	96.76	102.16	3,566	14,937
655	338,692	89,787	1,354,333	294,261	41.0	38.9	59.7	56.4	97.42	102.18	3,584	15,065
660	343,661	91,050	1,375,412	298,407	41.0	38.9	59.7	56.4	98.39	104.14	3,638	15,193
665	344,870	91,429	1,380,771	299,755	41.1	38.9	59.8	56.4	99.91	104.37	3,692	15,321
670	348,527	91,880	1,396,217	301,289	41.1	38.9	59.8	56.4	100.27	106.14	3,746	15,449
675	350,784	93,085	1,405,870	305,413	41.2	38.9	59.9	56.4	101.39	106.59	3,800	15,577
680	355,725	93,117	1,426,241	305,554	41.2	38.9	59.9	56.4	102.08	107.32	3,851	15,703
685	356,695	93,626	1,430,176	307,523	41.2	38.9	59.9	56.5	103.57	109.01	3,854	15,889
690	358,880	94,706	1,439,305	311,339	41.2	38.9	59.9	56.5	103.87	109.07	3,858	16,075
695	359,773	95,483	1,442,851	313,900	41.2	39.0	59.9	56.6	104.54	109.94	3,861	16,261
700	359,977	97,371	1,443,643	320,751	41.2	39.0	59.9	56.6	104.81	111.08	3,865	16,447
705	359,977	98,049	1,443,643	322,853	41.2	39.0	59.9	56.6	104.87	111.93	3,868	16,638

APPENDIX A. DETAILED RESULTS

Table A.20: Detailed results for Commons Collections (645-705 seconds)

	Ger Sequ	lerated		overed Inputs	Line Co	verage (%)	Method C	overage (%)	Cy. Complexit	Asserts/ clomatic y (6,848)	Ë.	illing Tests
Time	RANDOOP	RAIKON	RANDOOP	Raikon	RANDOOP	Raikon	Randoop	Raikon	RANDOOP	Raikon	RANDOOP	Raikon
710	360,104	98,447	1,444,321	326,451	41.2	39.0	59.9	56.6	104.87	114.50	3,871	16,764
715	360,105	99,073	1,444,321	327,007	41.2	39.0	60.0	56.7	104.91	115.46	3,875	16,890
720	360,105	99,660	1,444,321	328,496	41.2	39.0	60.0	56.7	104.91	116.79	3,878	17,016
725	360,105	100,845	1,444,321	332,705	41.2	39.0	60.0	56.7	104.91	117.66	3,882	17,142
730	360,108	101,362	1,444,322	334,519	41.2	39.0	60.0	56.7	104.91	119.04	3,885	17,264
735	360,120	101,698	1,444,368	335,632	41.2	39.1	60.0	56.8	104.91	119.73	3,888	17,425
740	360,120	103,072	1,444,368	340,553	41.2	39.1	60.0	56.8	104.91	120.20	3,892	17,586
745	360,120	104,455	1,444,368	345,565	41.2	39.2	60.0	56.8	104.91	121.96	3,895	17,747
750	360,123	104,620	1,444,378	346,094	41.2	39.2	60.0	56.8	104.91	123.62	3,899	17,908
755	360,199	106,281	1,444,671	352,095	41.2	39.3	60.0	56.8	104.91	123.96	3,902	18,079
760	360,223	106,281	1,444,766	352,095	41.2	39.3	60.0	56.8	104.94	125.00	3,905	18,198
765	360,290	106,303	1,445,035	352,180	41.2	39.4	60.0	56.9	104.94	125.95	3,909	18,317
770	360,515	106,700	1,446,041	356,387	41.2	39.4	60.0	56.9	104.96	125.95	3,912	18,436
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Table A.21: Detailed results for Commons Collections (710-770 seconds)

	Gen Seqt	erated Jences	_ 6	vered	Line Co	verage (%)	Method C	overage (%)	Cy Complexit	Asserts/ clomatic y (6,848)		ailing 「ests
Time	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon	Randoop	Raikon
775	360,515	108,133	1,446,041	358,428	41.2	39.5	60.0	57.0	105.03	125.96	3,916	18,555
780	360,831	108,503	1,447,210	359,735	41.2	39.5	60.0	57.0	105.03	128.54	3,922	18,670
785	360,836	108,550	1,447,210	359,903	41.2	39.5	60.0	57.0	105.12	129.14	3,923	18,707
790	360,839	109,240	1,447,283	362,176	41.2	39.5	60.0	57.0	105.13	129.21	3,924	18,744
795	360,839	109,612	1,447,314	363,497	41.2	39.5	60.0	57.1	105.13	130.14	3,924	18,781
800	360,891	110,713	1,447,515	367,619	41.2	39.5	60.0	57.1	105.13	130.54	3,925	18,818
805	360,903	110,936	1,447,515	368,415	41.2	39.5	60.0	57.1	105.14	131.88	3,926	18,863
810	360,903	111,280	1,447,638	369,579	41.2	39.5	60.0	57.1	105.15	132.23	3,927	18,986
815	360,951	112,501	1,447,710	371,586	41.2	39.5	60.0	57.1	105.15	132.63	3,928	19,109
820	360,952	113,016	1,447,710	375,511	41.2	39.6	60.0	57.2	105.16	134.88	3,928	19,232
825	361,379	113,285	1,449,262	376,460	41.2	39.6	60.0	57.2	105.16	135.35	3,929	19,355
830	361,379	113,435	1,449,262	376,891	41.2	39.6	60.0	57.2	105.30	135.61	3,930	19,478
835	361,670	114,222	1,450,554	379,712	41.2	39.6	60.0	57.2	105.30	136.53	3,931	19,613

Table A.22: Detailed results for Commons Collections (775-835 seconds)

APPENDIX A. DETAILED RESULTS

	Ger Seqı	nerated Jences	<u> </u>	vered	Line Co	overage (%)	Method C	overage (%)	Cyc	Asserts/ clomatic y (6,848)	Ë,	illing Tests
Time	Randoop	Raikon	RANDOOP	Raikon	RANDOOP	Raikon	Randoop	Raikon	RANDOOP	Raikon	RANDOOP	Raikon
840	361,670	114,269	1,450,554	379,915	41.2%	39.6%	60.0%	57.3%	105.38	136.61	3,932	19,748
845	361,691	115,614	1,450,625	384,830	41.2%	39.6%	60.0%	57.3%	105.38	138.43	3,932	19,883
850	361,765	116,329	1,450,912	387,399	41.2%	39.6%	60.0%	57.3%	105.39	139.38	3,933	20,018
855	361,777	117,493	1,450,912	391,724	41.2%	39.6%	60.0%	57.3%	105.40	140.12	3,934	20,149
860	361,777	117,935	1,451,327	393,503	41.2%	39.6%	60.0%	57.3%	105.41	140.86	3,935	20,306
865	361,796	118,729	1,451,331	397,548	41.2%	39.6%	60.0%	57.3%	105.41	141.39	3,936	20,463
870	361,866	119,616	1,451,348	399,334	41.2%	39.7%	60.0%	57.4%	105.41	143.38	3,936	20,620
875	361,895	121,166	1,451,407	404,758	41.2%	39.7%	60.0%	57.4%	105.44	145.27	3,937	20,777
880	361,942	122,394	1,451,657	409,183	41.2%	39.7%	60.0%	57.4%	105.44	146.61	3,938	20,930
885	361,942	122,536	1,451,657	409,771	41.2%	39.7%	60.0%	57.4%	105.46	146.81	3,938	21,007
890	361,942	122,796	1,451,657	410,652	41.2%	39.7%	60.0%	57.4%	105.46	147.10	3,941	21,033
895	362,033	122,913	1,452,247	411,058	41.2%	39.7%	60.0%	57.4%	105.48	147.25	3,941	20,964
006	362,076	122,940	1,452,275	411,131	41.9	39.7	60.0	57.4	105.50	147.41	3,940	21,040
			Table A.	23: Detaile	d results fc	or Commo	ins Collectic	ns (840-9(00 seconds)			

(840-900 seconds)
Collections
Commons
results for
.23: Detailed
le A

Appendix B

Additional Tooling

During this master thesis project some tools were created to aid in several aspects of the process. This chapter describes the developed tools and what they are meant for.

В.1 вів2сѕи

The first tool developed was bib2csv. This tool was developed in collaboration with Márcio Coelho and is meant to take a BibTeX file properly annotated with keywords, and converts it to a csv file compatible with the CoNExP tool [Yevtushenko, 2011]. This was done in order to facilitate the application of formal concept analysis in the context of literature revision. The tool is available as an open source Perl script at https://github.com/TiagoVeloso/Bib2CSV.

B.2 Fan-In/Fan-Out Calculator

This is another tool that was developed in the context of this master thesis was a simple Fan-In/Fan-Out Calculator generator. This tool was built on top of the already existing javap tool, which is a Java class disassembler. The javap tool is able to extract class dependencies. From this, all it was done was to apply a sequence of regular expressions to transform the output of javap into a useable format, dot or csv. It was written in Perl, the tool is available under an open source licence at https://github.com/TiagoVeloso/CallGraph.